

GMRES with embedded ensemble propagation for the efficient solution of parametric linear systems in uncertainty quantification of computational models with application to the thermomechanical simulation of an ITER front mirror





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Outline

(1) Introduction

- (2) Ensemble propagation
- (3) Ensemble parametric linear systems
- (4) Ensemble GMRES without ensemble reduction
- (5) Ensemble propagation of mesh-tying problems
- (6) Ensemble propagation of contact problems
- (7) Application to the front mirror of ITER
- (8) Conclusions and future works

Context: ITER front mirror

Multiphysics models





Context: ITER front mirror



Front mirrors of ITER diagnostic systems are subject to high thermal loads emitted by the plasma which thermally deform the mirror and can reduce the optical quality of the measures.

Multiphysics models

Context: Parametric computations

In order to quantify the robustness of the design to uncertainties in the parameters, the numerical model f of the front mirror must be evaluated a large number of times N:



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As N is typically large and f costly to evaluate, the total CPU cost of the parametric computation is large.

Context: High Performance Computing



High Performance Computing is an opportunity for parametric computation:

 Clusters will allow to run larger and larger number of evaluations concurrently,





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The high performance computing library Trilinos provides embedded capabilities targeting parametric computations.



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- Developed by 225 contributors.





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Trilinos provides a solver stack which includes definitions of matrices, linear solvers, nonlinear solvers, preconditioners, ... templated on the data type.





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Trilinos provides a solver stack which includes definitions of matrices, linear solvers, nonlinear solvers, preconditioners, ... templated on the data type.





This feature is a key feature for the work presented today.

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- ▶ Is a C++ Performance portability library;
- Enables single source performance portable codes;



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Those **data abstractions** and **programming models** are used in the work presented today to implement the **efficient GEMV** discussed later.



Stokhos:

 Is the Trilinos package for embedded uncertainty quantification methods;



Parametric computations

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- Provides an implementation of the embedded ensemble propagation method;
- Is developed by Eric T. Phipps.



The work presented today uses and contributes to the embedded ensemble propagation method.

Objective of the thesis

For a given set of samples, to **reduce** the **wall-clock time** to evaluate multiphysics models on high performance clusters.



This has been done using and contributing to the **embedded ensemble propagation** one of the embedded strategies implemented in **Stokhos**.

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In sampling-based parametric computation, instead of individually evaluating each instance of the model, Ensemble propagation (EP) consists of **simultaneously evaluating** a **subset of samples** of the model.

Given N samples and an ensemble size s, instead of looping over the N samples:

```
for i in range(0,N):
    y[i] = f(p[i])
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for i in range(0,N):
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```



Ensemble propagation loops over sets of size s of the N samples:

```
for i in range(0,N,s):

y[i:i+s] = f(p[i:i+s])

p^{(i)}, \ldots, p^{(i+s-1)}
f
y^{(i)}, \ldots, y^{(i+s-1)}
```

Using EP increases the order of sample-dependent tensors by one:

Without EP With EP



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Algorithms should be adapted accordingly.

EP was introduced by [Phipps, 2017], made available in **Stokhos**, and implemented using a **template-based generic-programming** approach:

```
template <typename T, int ensemble_size>
class Ensemble{
  T data[ensemble_size];
  Ensemble<T,ensemble_size> operator+ (const Ensemble<T,ensemble_size> &v);
  Ensemble<T,ensemble_size> operator- (const Ensemble<T,ensemble_size> &v);
  Ensemble<T,ensemble_size> operator* (const Ensemble<T,ensemble_size> &v);
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and providing template specializations for some of the Trilinos functions and classes for this new data type.

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```

and providing template specializations for some of the Trilinos functions and classes for this new data type.

This implementation strategy allows to use EP in the full solver stack of **Trilinos** supporting templated data types.

Advantages of the EP:

- Reuse of common variables;
- More opportunities for vectorization (more data parallelism);
- Improved memory access pattern;
- ▶ Reduction of Message Passing Interface (MPI) latency per sample.

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Challenges of the EP:

- Increased memory usage compared to a single computation;
- Ensemble divergence:
 - control flow divergence:
 - If-then-else divergence;
 - Loop divergence;
 - function call divergence:
 - Missing BLAS functions;
 - What is an inner product of two ensemble typed vectors?

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Parametric linear systems

We want to solve a **parametric linear system** for a subset of *s* samples of the parameters together:

$$\boldsymbol{A}_{::\ell} \, \boldsymbol{x}_{:\ell} = \boldsymbol{b}_{:\ell}$$
 for all $\ell = 1, \dots, s$,

where matrices $A_{::1}, \ldots, A_{::s}$ are sparse and not symmetric or not positive definite. Representation of a system for s = 4:



How to solve the sparse parametric linear system efficiently with EP?
Right-preconditioned GMRES

 $r^{(0)} = b - A x^{(0)}$ $\beta = \|\boldsymbol{r}^{(0)}\|$ $v_{11} = r^{(0)}/\beta$ for j = 1, ..., m do $\mathbf{w} = \mathbf{A}\mathbf{M}^{-1}\mathbf{v}_{\cdot i}$ $\boldsymbol{h}_{(1:j)j} = \boldsymbol{V}_{:(1:j)}^{\mathrm{T}} \boldsymbol{w}$ $\mathbf{v}_{:(i+1)} = \mathbf{w} - \mathbf{V}_{:(1:j)} \mathbf{h}_{(1:j)j}$ $h_{(i+1)i} = \| \mathbf{v}_{i,(i+1)} \|$ if $h_{(i+1)i} \neq 0$ then $| \mathbf{v}_{(i+1)} = \mathbf{v}_{(i+1)} / h_{(i+1)}$ else m = i_ break if $oldsymbol{q}_{:(j+1)}^{\mathrm{T}}oldsymbol{e}_1 \leq arepsilon$ then m = ibreak $\mathbf{y} = \arg\min_{\mathbf{z}} \| \beta \, \mathbf{e}_1 - \mathbf{H}_{(1:m+1)(1:m)} \, \mathbf{z} \|$ $x^{(m)} = x^{(0)} + M^{-1} V_{:(1:m)} y$

Right-preconditioned GMRES

 $r^{(0)} = b - A x^{(0)}$ $\beta = \|\boldsymbol{r}^{(0)}\|$ $v_{:1} = r^{(0)} / \beta$ for j = 1, ..., m do $\mathbf{w} = \mathbf{A}\mathbf{M}^{-1}\mathbf{v}_{\cdot i}$ $oldsymbol{h}_{(1:j)j} = oldsymbol{V}_{:(1:j)}^{ ext{T}}oldsymbol{w}$ $\mathbf{v}_{:(i+1)} = \mathbf{w} - \mathbf{V}_{:(1:i)} \mathbf{h}_{(1:i)i}$ $h_{(j+1)j} = \| \mathbf{v}_{(j+1)} \|$ if $h_{(i+1)i} \neq 0$ then $\mathbf{v}_{(i+1)} = \mathbf{v}_{(i+1)}/h_{(i+1)}$ else m = ibreak if $\boldsymbol{q}_{:(i+1)}^{\mathrm{T}} \boldsymbol{e}_1 \leq \varepsilon$ then m = ibreak $\mathbf{y} = \arg\min_{\mathbf{z}} \|\beta \, \mathbf{e}_1 - \mathbf{H}_{(1:m+1)(1:m)} \, \mathbf{z}\|$ $x^{(m)} = x^{(0)} + M^{-1} V_{(1:m)} y$

The GMRES method iteratively creates an orthonormal basis $v_{:1}, \ldots, v_{:j}$ for the vector space:

$$\operatorname{span}\left\{ \pmb{r^{(0)}}, \pmb{AM^{-1}r^{(0)}}, \dots, \left(\pmb{AM^{-1}} \right)^{(j-1)} \pmb{r^{(0)}}
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using the orthonarmal basis of the current iteration $\mathbf{v}_{:1}, \ldots, \mathbf{v}_{:j}$, applying \mathbf{AM}^{-1} to the last vector $\mathbf{v}_{:j}$ and orthonormalize it with $\mathbf{v}_{:1}, \ldots, \mathbf{v}_{:j}$.

This process continues up to the point where the basis allows to have a sufficiently small error to the solution of the linear system.

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This process continues up to the point where the basis allows to have a sufficiently small error to the solution of the linear system.

What is an inner product of two ensemble typed vectors?

Reduced inner product

First approach [Phipps, 2017]: equivalent to gathering the sample matrices $A_{::1}, \ldots, A_{::s}$ into a block diagonal matrix

$$\begin{bmatrix} \mathbf{A}_{::1} & & \\ & \ddots & \\ & & \mathbf{A}_{::s} \end{bmatrix} \begin{bmatrix} \mathbf{x}_{:1} \\ \vdots \\ \mathbf{x}_{:s} \end{bmatrix} = \begin{bmatrix} \mathbf{b}_{:1} \\ \vdots \\ \mathbf{b}_{:s} \end{bmatrix},$$

and to applying an iterative method on the block diagonal system.

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$$egin{bmatrix} egin{aligned} egi$$

and to applying an iterative method on the block diagonal system.

This is mathematically equivalent to defining a reduced inner product:

Advantages: No ensemble divergence and possibility to use efficient BLAS implementations, **Challenges:** The samples are coupled together, the spectra of $A_{::1}, \ldots, A_{::s}$ are gathered, and the condition number increases and is larger than the ones of $A_{::1}, \ldots, A_{::s}$. The number of iterations increases.

Ensemble-typed inner product

Second approach [D'Elia, 2020]: to avoid the coupling of the samples together using an **ensemble-typed inner product**:



It was first introduced for grouping purpose.

Ensemble-typed inner product

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It was first introduced for grouping purpose.

Advantage: No coupling: each sample converges as fast as if it was propagated alone, **Challenge:** Every ensemble divergence has to be managed explicitly.

With ensemble reduction:

Advantages:

- ► No control flow divergence;
- Use of standard libraries such as the Intel Math Kernel Library.

Challenges:

- Convergence in the least-squares sense;
- The spectrum of the ensemble matrix is the union of the spectra of the sample matrices: having a good preconditioner is more complex;
- Increased number of iterations.

Advantages and challenges of both approaches

With ensemble reduction:

Advantages:

- ► No control flow divergence;
- Use of standard libraries such as the Intel Math Kernel Library.

Challenges:

- Convergence in the least-squares sense;
- The spectrum of the ensemble matrix is the union of the spectra of the sample matrices: having a good preconditioner is more complex;
- Increased number of iterations.

Without ensemble reduction:

Advantages:

- Convergence of an iterative method implies the convergence for every sample;
- ► The spectra are not gathered;
- Global convergence rate controlled by the slowest sample.

Challenges:

- Control flow divergence has to be treated explicitly;
- No current implementation of the needed BLAS routines in the Intel Math Kernel Library.

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Occurrence of ensemble divergence in GMRES

 $r^{(0)} = h - A x^{(0)}$ $\beta = \|\boldsymbol{r}^{(0)}\|$ $v_{11} = r^{(0)}/\beta$ for i = 1, ..., m do $\boldsymbol{w} = \boldsymbol{A}\boldsymbol{M}^{-1}\,\boldsymbol{v}_{:i}$ $oldsymbol{h}_{(1:j)j} = oldsymbol{V}_{:(1:j)}^{ ext{T}}oldsymbol{w}$ $\mathbf{v}_{:(i+1)} = \mathbf{w} - \mathbf{V}_{:(1:i)} \mathbf{h}_{(1:i)i}$ $h_{(i+1)|i} = \|\mathbf{v}_{i,(i+1)}\|^2$ if $h_{(j+1),i} \neq 0$ then $| \mathbf{v}_{:(j+1)} = \mathbf{v}_{:(j+1)} / h_{(j+1)j}$ else m = ibreak if $\boldsymbol{q}_{:(i+1)}^{\mathrm{T}} \boldsymbol{e}_1 \leq \varepsilon$ then m = ibreak

$$y = \arg \min_{z} \|\beta e_{1} - H_{(1:m+1)(1:m)} z\|$$
$$x^{(m)} = x^{(0)} + M^{-1} V_{:(1:m)} y$$

Ensemble divergence in GMRES:

1. a vector can require a normalization or not: **if-then-else divergence**;

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Ensemble divergence in GMRES:

- 1. a vector can require a normalization or not: **if-then-else divergence**;
- different samples may require different numbers of iterations to converge: loop divergence;

$$\begin{split} & \mathbf{y} = \arg\min_{\mathbf{z}} \|\beta \ \mathbf{e}_1 - \mathbf{H}_{(1:m+1)(1:m)} \ \mathbf{z} \| \\ & \mathbf{x}^{(m)} = \mathbf{x}^{(0)} + \mathbf{M}^{-1} \mathbf{V}_{:(1:m)} \ \mathbf{y} \end{split}$$

```
typedef EnsembleTrait<T> ET;
const int s = ET::ensemble_size;
bool all_zeros = true;
```

```
typedef EnsembleTrait<T> ET;
const int s = ET::ensemble_size;
bool all_zeros = true;
T norm_inv;
for (int l = 0; l < s; ++1)
if (ET::coeff(norm, l) > 0) {
ET::coeff(norm_inv, l) = 1. / ET::coeff(norm, l);
all_zeros = false;
}
else
ET::coeff(norm_inv, l) = 0.;
```

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typedef EnsembleTrait<T> ET;
const int s = ET::ensemble_size;
bool all_zeros = true;
T norm inv:
for (int l = 0; l < s; ++1)
  if (ET::coeff(norm, 1) > 0) {
    ET::coeff(norm_inv, l) = 1. / ET::coeff(norm, l);
    all_zeros = false;
  3
  else
    ET::coeff(norm inv, 1) = 0.:
if (all_zeros) return has_converged;
for (int i = 0; i < n; ++i)</pre>
  v[i] *= norm_inv;
```

```
typedef EnsembleTrait<T> ET;
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  for (int i = 0; i < n; ++i)</pre>
    v[i] *= norm_inv;
This is correct, but not concise and relies on the compiler to optimize the code.
```

Control flow divergence: Mask class

The control flow divergence, both the **if-then-else divergence** and the **loop divergence**, can be solved by defining a **Mask class** equivalent to:

```
template <int ensemble_size>
class Mask{
   bool data[ensemble_size];
   //...
}
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which stores the result of any comparison of ensembles sample-wise.

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This mask can then used for masked assignments and logical reductions:

```
T norm_inv;
if (AND(norm == 0)) return has_converged;
MaskAssign(norm > 0, norm_inv) /= {1., norm, 0.};
for (int i = 0; i < n; ++i)
v[i] *= norm inv;
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v[i] *= norm_inv;
```

This second implementation is **more concise**, potentially **more readable**, and **helps** the optimization performed by the compiler.

Occurrence of ensemble divergence in GMRES

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$$\begin{aligned} & \mathbf{y} = \arg\min_{\mathbf{z}} \|\beta \, \mathbf{e}_1 - \mathbf{H}_{(1:m+1)(1:m)} \, \mathbf{z}\| \\ & \mathbf{x}^{(m)} = \mathbf{x}^{(0)} + \mathbf{M}^{-1} \mathbf{V}_{:(1:m)} \, \mathbf{y} \end{aligned}$$

Ensemble divergence in GMRES:

- 1. a vector can require a normalization or not: **if-then-else divergence**;
- different samples may require different numbers of iterations to converge: loop divergence;
- called BLAS functions, such as GEMV for the dense matrix-vector operations in the orthogonalization process (inner products and update), may not support ensemble-typed inputs, leading to function call divergence.

The **GEMV** with ensemble propagation takes the form of **tensors contractions** as follows:

$$\mathbf{y}_{:\ell} = \beta_{\ell} \, \mathbf{y}_{:\ell} + \alpha_{\ell} \, \mathbf{A}_{::\ell} \, \mathbf{x}_{:\ell} \quad \text{for all} \quad \ell = 1, \dots, s,$$

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Interleaved memory layout of the $n \times j \times s$ third-order tensor \mathcal{A} due to EP:

$$a_{ik\ell} \leftrightarrow a\left[(i-1)s+(k-1)ns+(\ell-1)
ight],$$

with n the number of degrees of freedom per sample, j the Krylov subspace dimension, and s the ensemble size;

Tall skinny third-order tensor ${\cal A}$



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$$\mathbf{y}_{:\ell} = \beta_\ell \, \mathbf{y}_{:\ell} + \alpha_\ell \, \mathbf{A}_{::\ell} \, \mathbf{x}_{:\ell} \quad \text{for all} \quad \ell = 1, \dots, s,$$

Interleaved memory layout of the $n \times j \times s$ third-order tensor \mathcal{A} due to EP:

$$a_{ik\ell} \leftrightarrow a\left[(i-1)s+(k-1)ns+(\ell-1)
ight],$$

with n the number of degrees of freedom per sample, j the Krylov subspace dimension, and s the ensemble size;

Challenge: the **memory layout** prevents us from using efficiently a **scalar-typed GEMV** implementation sequentially *s* times.

Tall skinny third-order tensor \mathcal{A}



Such an operation has a **low arithmetic intensity** as, for every $a_{ik\ell}$ loaded from memory only two operations are performed.

The throughput of this computation is therefore limited by the **memory bandwidth** on standard architectures. The speed-up of this tensors contraction versus *s* GEMV with unit stride **cannot be greater than 1** providing that both implementation reach maximal throughput.

How should we implement the contraction such that theoretical performance is achieved?

To reach full bandwidth, we have to:

- **Exploit the parallelism of the architecture:**
 - ▶ Use every physical core as much as possible.

► Transfer data efficiently through the memory hierarchy:

- Keep reusable data in cache;
- Use unit stride loads.

► Exploit CPU power:

- Keep reusable data in registers;
- ▶ Use vector load and store, avoid vector gather.

S

parfor
$$t = 1$$
 to $n - n_c + 1$ by m_c do
for $i = t, ..., t + n_c - 1$ do
 $\lfloor y_{i\ell} = \beta_\ell y_{i\ell}$ for all $\ell = 1, ..., s$
for $k = 1, ..., j$ do
 $\gamma_\ell = \alpha_\ell x_{k\ell}$ for all $\ell = 1, ..., s$
for $i = t, ..., t + n_c - 1$ do
 $\lfloor y_{i\ell} = y_{i\ell} + \gamma_\ell a_{ik\ell}$ for all $\ell = 1, ..., s$

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28 / 55

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 n_c n Tile level

Outer level

28 / 55

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Tiling:

- Each thread applies a tile at a time;
- ► Cache blocking of **Y**.
- Vectorization:
 - Vectorization of the loops over the samples;
 - Intel Intrinsics or overloaded operators.



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Tiling:

- Each thread applies a tile at a time;
- ► Cache blocking of **Y**.
- Vectorization:
 - Vectorization of the loops over the samples;
 - Intel Intrinsics or overloaded operators.
- Choice of the tile size n_c to keep $\mathbf{Y}_{(t:t+n_c-1)}$: in cache.



Outer level

Choice of the tile size n_c : example on Intel 8100-Series (Skylake)

Measured bandwidth (1 NUMA region): 101.2108 GB/s Deduced maximal throughput: 25 GFLOPS Parameters:

- ► Third order tensor A of size n × j × s;
- n = 768000, j = 300, s is the ensemble size;
- only evaluated n_c such that n is a multiple of Nn_c where N = 48 is the number of threads.

The highlighted value of $n_c^{\max}s =$ 32768 corresponds to the case where $\mathbf{Y}_{(t:t+n_c-1):}$ takes at most half of the L2 cache [Goto, 2008].

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GEMV: results - Intel(R) Xeon(R) Platinum 8160 CPU

The tile size is chosen based on n and n_c^{\max} to have evenly distribute work among the threads:

$$n_c = \left\lceil \frac{n}{N \left\lceil \frac{n}{N n_c^{\text{max}}} \right\rceil} \right\rceil.$$

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Parameters:

- Threads N = 48;
- ▶ n = 30000;
- n_c^{max} = 4096, 2048, 1365, and 1024 for s = 8, 16, 24, and 32 respectively.

Performance similar to the ensemble reduction with MKL used to compute the update.



In [Liegeois, 2020], we describe and implement an efficient ensemble GMRES without ensemble reduction.

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The control flow divergence and the function call divergence have been solved by:

- Implementing a Mask class which is used to apply masked assignment and logical reduction;
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The control flow divergence and the function call divergence have been solved by:

- Implementing a Mask class which is used to apply masked assignment and logical reduction;
- ► Implementing an efficient ensemble GEMV for the orthogonalization process.

Those two contributions lead to:

- ► An equivalent cost per iteration of ensemble GMRES with and without reduction;
- A safe implementation which is able to deal with early-converged samples.

Example using a discretized Dirichlet problem for the 1D Laplacian with n = 6:

$$\mathbf{A}\mathbf{x} = \mathbf{b},$$
where:

$$\mathbf{A} = \kappa \begin{bmatrix} 2 & -1 & & \\ -1 & 2 & -1 & & \\ & \ddots & \ddots & \ddots & \\ & & -1 & 2 & -1 \\ & & & -1 & 2 \end{bmatrix}.$$

Samples are considered:

1.
$$\kappa = 1$$
, $\boldsymbol{b}^{\mathrm{T}} = \begin{bmatrix} 0 & 0 & 0 & 1 & 0 & 0 \end{bmatrix}$,
2. $\kappa = 1$, $\boldsymbol{b}^{\mathrm{T}} = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$,
3. $\kappa = 1.5$, $\boldsymbol{b}^{\mathrm{T}} = \begin{bmatrix} 0 & 0 & 0 & 1 & 0 & 0 \end{bmatrix}$



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Coupling 1 and 2 changes the convergence but not the number of iterations,

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Coupling 1 and 3 multiply by 2 the number of iterations.

Katoptron

- Fully templated C++ code heavily based on Trilinos which provides a fully templated solver stack;
- Embedded in a Python interface. This eases the looping around samples, the grouping of samples together, etc;
- Hybrid parallelism based on Tpetra with MPI for distributed memory and Kokkos with OpenMP for shared memory;
- Uses Gmsh to import 3D meshes and VTK to write the output files;
- Open source and freely available on https://gitlab.uliege.be/am-dept/waves.



Outline

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- (6) Ensemble propagation of contact problems
- (7) Application to the front mirror of ITER
- (8) Conclusions and future works

Mesh-tying problem

- Plate with a hole pulled on two opposite sides;
- Two meshes glued with the Mortar finite element method in saddle point formulation;
- Lamé parameters represented as a lognormal random field;
- Multigrid preconditioner with saddle point matrix on each multigrid level;
- 32088 degrees of freedom per sample;
- ▶ 1 Intel Skylake CPU.



Shear modulus [GPa]

Preconditioners: Full multigrid approach

Introduced in [Wiesner, 2015] for contact problem.

- Main idea: use coarser representations of fine level problems in order to speed up the solution process;
- Uses the multigrid approach on the full matrix, preserving the saddle-point structure on all levels;
- Algebraic multigrid: no special information is necessary to build the multigrid hierarchies;
- Mutligrid hierarchies are independent of the activity of the Lagrange multipliers;
- Implemented in MueLu with already existing codes and contributions of this thesis;
- Due to EP, level matrices are now third-order tensors but prolongation and restriction operators are sample independent.



Mesh-tying example: speed-up of one GMRES iteration

Speed-up: relative gain in wall-clock time (architecture dependent):

$$\mathsf{S}(e) = rac{\sum_{\ell \in e} \mathsf{Time}_{\ell}}{\mathsf{Time}_{e}}.$$

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The speed-up per iteration is nearly independent of the use of ensemble reduction.













The reduced number of iterations to converge without ensemble reduction improves the speed-up compared to ensemble reduction.

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Beam contact problem

- Size: L = 50 cm, W = 5 cm, H = 5 cm, d = 1 cm;
- Elements: $60 \times 6 \times 6$ hexahedra;
- Number of Dofs: 9394;
- Depending on the pressure p, the contact is fully open or partially closed;
- Uncertainties:
 - > Young's modulus: $E \sim \mathcal{U}(205, 215)$ [GPa];
 - ▶ Pressure: $p \sim U(5, 25)$ [MPa].
- Quantity of Interest: displacement along z at point (L, 0, H/2);
- 640 Halton Quasi Monte Carlo samples;
- Solved on 1 Intel Skylake CPU.



Ensemble propagation for mechanical contact problem

Instead of individually solving the mechanical contact problem for each instance of the model, we have to **solve simultaneously** the mechanical contact problem for **a subset of samples** of the model.

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Challenges of the EP for mechanical contact problem:

▶ Different samples can have different active Lagrange multipliers,



Samples may require a different number of active set iterations.

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Challenges of the EP for mechanical contact problem:

► Different samples can have different active Lagrange multipliers,



Samples may require a different number of active set iterations.

Those challenges have been solved using the developed **Mask class** and waiting for the **convergence of all samples** in the active set strategy.

Contact example: convergence and total speed-up



Contact example: convergence and total speed-up



Contact example: convergence and total speed-up



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Front mirror context

First mirrors of optical diagnostics in ITER:

- are exposed to high radiation and fluxes of particles which escape the plasma;
- are the most vulnerable in-vessel optical components, being subject to erosion or to deposition of impurities.

Material selected for the reflecting surface must combine:

- a high optical reflectivity in a wide spectral range;
- > a sufficient resistance to physical sputtering.

Rhodium^a is identified as a promising candidate, due to:

- low sputtering in most cases;
- high optical reflectance;
- optical reflectance insensitive to large temperature changes.



^aP. Mertens, R. Boman, S. Dickheuer, Y. Krasikov, A. Krimmer, D. Leichtle, K. Liegeois, C. Linsmeier, A. Litnovsky, O. Marchuk, M. Rasinski, M. De Bock, On the use of rhodium mirrors for optical diagnostics in ITER, *Fusion Engineering and Design*, 146:2514–2518, 2019.

ITER test case

- Thermomechanical problem;
- The contact interfaces are not modeled, the mesh is fused at the common interface of the components;
- The assembly is heated by surface and volumetric loads;
- Rigid body motions are prevented by setting zero displacements close to the bolt holes;
- Temperature at the cooling channel is set to 70 °C;
- 1.7 10⁶ elements and 1.31 10⁶ degrees of freedom per sample;
- 3-level multigrid preconditioner with block Gauss-Seidel level smoother and Klu as smoother on the coarsest level;
- Solved on 4 Intel Skylake CPUs.



ITER test case: Uncertainties and quantities of interest

- Uncertainties: the irradiation due to neutron damage of the mirror modifies the material properties of the used material; we consider the heat conductivity of the AIN (Aluminum nitride) ceramic spacers as random variables; k₁ is the thermal conductivity of the two small AIN spacers and k₂ is the thermal conductivity of the largest AIN spacer.
- Quantities of interest:
 - maximal temperature reached on the mirror surface;
 - deformation of the mirror surface;
- Uncertainty quantification strategy: evaluate the model at some samples to build a surrogate model using the nonintrusive spectral projection method.



ITER test case: Quantities of Interest



As the thermal conductivities influence the temperature distribution, they impact the deformation of the mirror surface and its curvature.
ITER test case: Surrogate models



As the thermal conductivities influence the temperature distribution, they impact the deformation of the mirror surface and its curvature.

ITER test case: total speed-up of ensemble GMRES without reduction



The total speed-up of ensemble GMRES is between 1.5 and 2 and is smaller than the speed-up of the sparse matrix vector product or the preconditioner due to the othogonalization process.

ITER test case: full assembly

- Thermomechanical problem on the full assembly;
- The contact interfaces are not modeled;
- 12.810⁶ elements and 9.1510⁶
 degrees of freedom per sample;
- 3-level multigrid preconditioner with block Gauss-Seidel level smoother and Klu as smoother on the coarsest level;
- Solved on 32 Intel Skylake CPUs.

5

Temperature



ITER test case: full assembly



The total speed-up of ensemble GMRES is about 1.5 and is smaller than the speed-up of the preconditioner due to the othogonalization process.

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Extension of the embedded ensemble propagation method from CG to GMRES:

Two variants of ensemble GMRES can currently be used: with ensemble reduction and without ensemble reduction;

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- Implementation of an efficient ensemble GEMV;
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- Implementation of a Mask class for ensemble types which is now used in ensemble GMRES with ensemble reduction;
- The implementation related to GMRES without reduction has been merged into the official Trilinos repository on github.

► Handling of contact constraints with embedded ensemble propagation;

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> An open-source software, Katoptron, has been developed to test ensemble GMRES;

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> Application to thermomechanical problems relevant for the **front mirrors** in ITER;

Observed total speed-up of about 2 in all cases.

Directions for future work

- Investigate the impact of the preconditioners on the influence of the reduction on the convergence;
- Investigate and apply the strategy on the mirror problem with contact losses;
- Investigate the use of ensemble GMRES on other problems such as fluid mechanics problems or ice sheet problems;
- Study grouping strategies for non-linear problems;
- Consider adaptive ensemble sizes;
- ► Test the method on transient problems.

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Contributions to the Trilinos repository

- Implemented from scratch:
 - Mask class, mask assignments, logical reduction in Stokhos;
 - Ensemble GEMV;
 - Coordinate type in Teuchos;
 - Map from Lagrange multipliers to DOF, and relative checks in Moertel;
 - ▶ Interface aggregation strategy, Interface mapping transfer.
- Implemented relying on specializations:
 - Ensemble GMRES without reduction;
 - Orthogonalization manager: DGKS, ICGS, IMGS;
 - Status tests: explicit and implicit;
 - ROTG, TRSM.
- Corrections of small bugs;
- Inclusion of new tests;
- Update the Trilinos automatic tests to compile EP with and without reduction;
- Packages impacted: Stokhos, Belos, Ifpack2, KokkosKernels, MueLu, Teuchos, Moertel, Xpetra, Galeri.

Mechanical contact problem



 $k \leftarrow 0$ Choose an initial guess for the active set \mathcal{A}_k **do** | Given \mathcal{A}_k , compute the solution of



Inner nodes: i, potential contact nodes: c, at iteration k, inactive set: \mathcal{I}_k , and active set: \mathcal{A}_k .

The algebraic full form as a way to handle activities

The matrix of the system:





- has a constant size but its graph varies with the active set,
- can be stored using an extended graph which is the union of all the possible graphs,
- has a saddle-point structure,
- ▶ is **not positive definite** (if at least one Lagrange multiplier is active).

Example sparse matrix vector product:

```
// CRS matrix-vector product z = A * x for arbitrary
// floating-point type T
template <typename T>
void crs mat vec(const CrsMatrix<T>& A.
                 const T *x, T *z) {
 for (int row =0: row<A.num rows: ++row) {</pre>
    const int entry_begin = A.row_map[row];
    const int entry end = A.row map[row+1];
   T \, sum = 0.0;
    for (int e = entry_begin; e<entry_end; ++e) {</pre>
      const int col = A.col_entry[e];
      sum += A.values[e] * x[col]:
    z[row] = sum:
```



GEMV: results - Intel(R) Xeon(R) Phi Knights Landing (KNL)

Xeon Phi KNL in quadrant cache mode Measured bandwidth of the MCDRAM: 320 GB/s

Deduced maximal throughput: 80 GFLOPS

Parameters:

- > Threads N = 128
- ▶ $n_c = 1024$ for s = 8, $n = 8 N n_c$,
- for a given n, data size independent of s.

Performance greater than the MKL,

Performance similar to the theoretical limit,

Sensibility to the order of the operations.





- The atomic adds introduced a fixed cost linked to the desynchronization of the threads that all want to access the first entries of the left-hand side vector at the same time.
- ▶ We used a cycling technique such that the threads start at different rows evenly distributed among *m*. This reduces the desynchronization cost for larger *m*.
- ▶ To reduce the fixed cost for small *m*, we gather threads per team of 4, do a parallel reduction per team and then do the atomics.