

## Extensions of the Pascal triangle to words

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# Classical Pascal triangle

$$P: (m, k) \in \mathbb{N} \times \mathbb{N} \mapsto \binom{m}{k} \in \mathbb{N}$$

$\binom{m}{k}$	$k$								
	0	1	2	3	4	5	6	7	...
0	1	0	0	0	0	0	0	0	
1	1	1	0	0	0	0	0	0	
2	1	2	1	0	0	0	0	0	
$m$	3	1	3	3	1	0	0	0	0
4	1	4	6	4	1	0	0	0	
5	1	5	10	10	5	1	0	0	
6	1	6	15	20	15	6	1	0	
7	1	7	21	35	35	21	7	1	
$\vdots$									$\ddots$

Binomial coefficients:

$$\binom{m}{k} = \frac{m!}{(m-k)!k!}$$

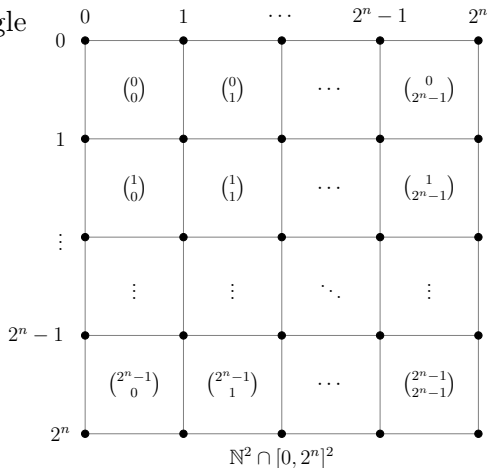
Pascal's rule:

$$\binom{m}{k} = \binom{m-1}{k} + \binom{m-1}{k-1}$$

# A specific construction

- Grid: first  $2^n$  rows and columns of the Pascal triangle

$$\left( \binom{m}{k} \bmod 2 \right)_{0 \leq m, k < 2^n}$$

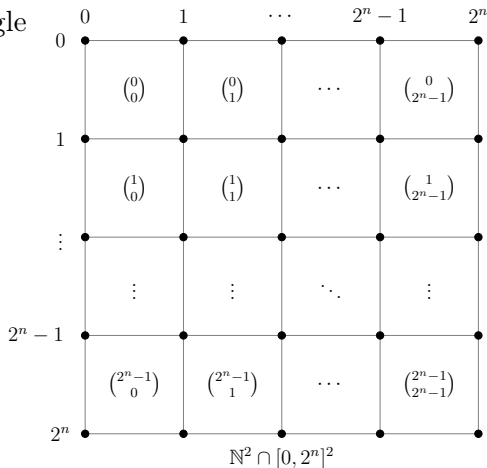


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  - white if  $\binom{m}{k} \equiv 0 \pmod 2$
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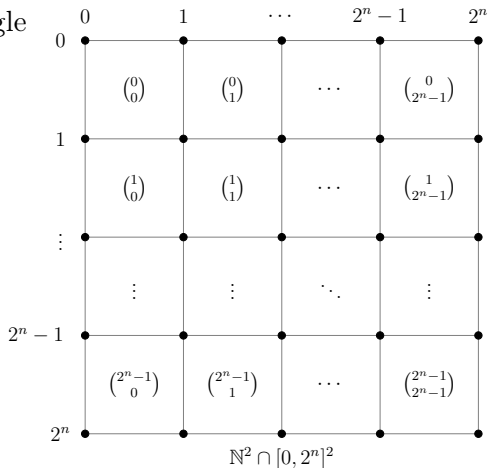


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- Normalize by a homothety of ratio  $1/2^n$  (bring into  $[0, 1]^2$ )



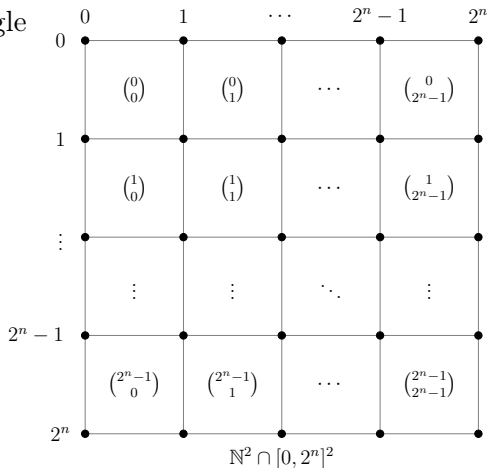
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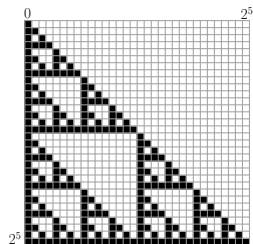
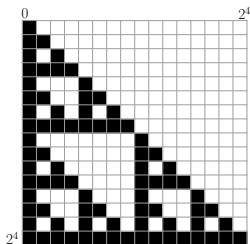
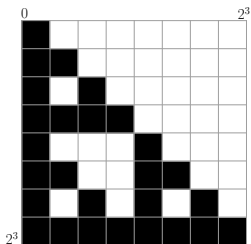
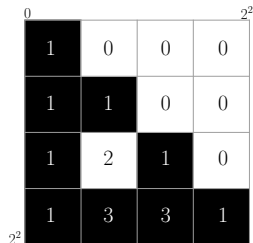
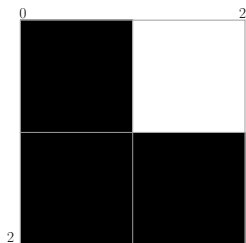
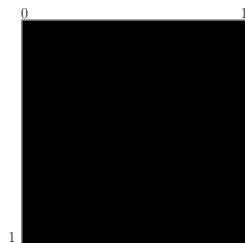
$$\left( \binom{m}{k} \bmod 2 \right)_{0 \leq m, k < 2^n}$$

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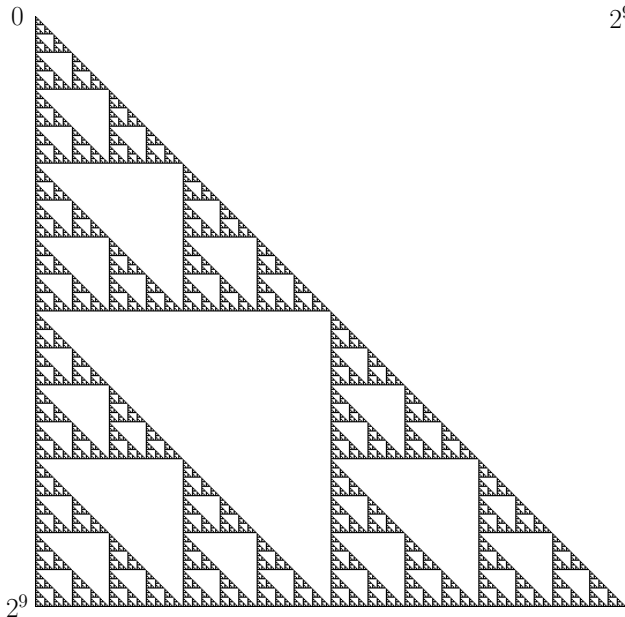
$\rightsquigarrow$  sequence of compact sets belonging to  $[0, 1]^2$



# The first six elements of the sequence



# The tenth element of the sequence





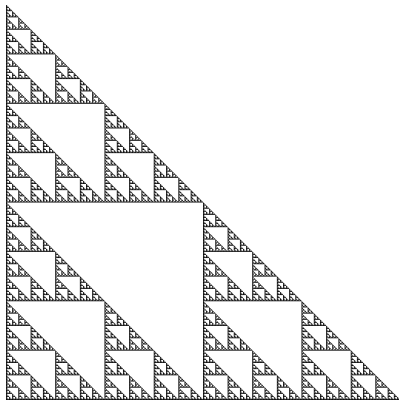
# The Sierpiński gasket



# The Sierpiński gasket

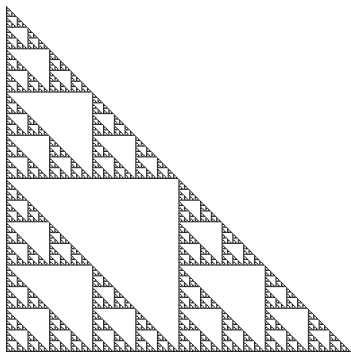


# The Sierpiński gasket



## Folklore fact

The latter sequence of compact sets converges to the Sierpiński gasket (w.r.t. the Hausdorff distance).



## Definitions:

- $\epsilon$ -fattening of a subset  $S \subset \mathbb{R}^2$

$$[S]_\epsilon = \bigcup_{x \in S} B(x, \epsilon)$$

- $(\mathcal{H}(\mathbb{R}^2), d_h)$  complete space of the non-empty compact subsets of  $\mathbb{R}^2$  equipped with the Hausdorff distance  $d_h$

$$d_h(S, S') = \inf\{\epsilon \in \mathbb{R}_{>0} \mid S \subset [S']_\epsilon \text{ and } S' \subset [S]_\epsilon\}$$

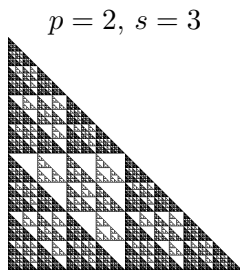
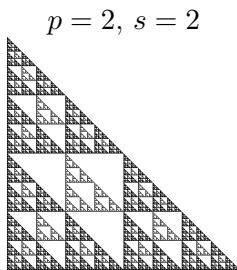
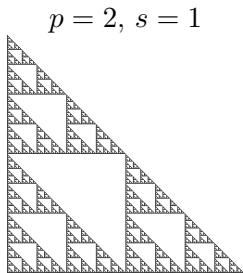
## Theorem (von Haeseler, Peitgen, and Skordev, 1992)

Let  $p$  be a prime and  $s > 0$ .

The sequence of compact sets corresponding to

$$\left( \binom{m}{k} \bmod p^s \right)_{0 \leq m, k < p^n}$$

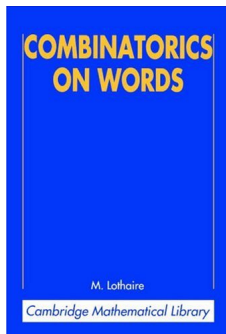
converges when  $n$  tends to infinity (w.r.t. the Hausdorff distance).



Replace integers by **finite words**.

## Combinatorics on words (CoW)

- new area of discrete mathematics ( $\pm 1900$ )
- study sequences of symbols (called letters)
- topics include:
  - ◇ regularities and patterns in words
  - ◇ important types of words (e.g. automatic, regular, de Bruijn, Lyndon, Sturmian)
  - ◇ coding of structures (e.g. paths, trees or curves in the plane)



M. Lothaire, 1983.

Definition: A *finite word* is a finite sequence of letters belonging to a finite set called *alphabet*.

Example:  $101, 101001 \in \{0, 1\}^*$

## Binomial coefficient of words

Let  $u, v$  be two finite words.

The *binomial coefficient*  $\binom{u}{v}$  of  $u$  and  $v$  is the number of times  $v$  occurs as a subsequence of  $u$  (meaning as a “scattered” subword).

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Example:  $u = 101001$        $v = 101$



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Example:  $u = 101001$        $v = 101$       1 occurrence

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Example:  $u = 101001$        $v = 101$       2 occurrences

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The *binomial coefficient*  $\binom{u}{v}$  of  $u$  and  $v$  is the number of times  $v$  occurs as a subsequence of  $u$  (meaning as a “scattered” subword).

Example:  $u = 101001$        $v = 101$       3 occurrences

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Example:  $101, 101001 \in \{0, 1\}^*$

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Example:  $u = 101001$        $v = 101$       4 occurrences

Definition: A *finite word* is a finite sequence of letters belonging to a finite set called *alphabet*.

Example:  $101, 101001 \in \{0, 1\}^*$

## Binomial coefficient of words

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Example:  $u = 10\mathbf{1001}$        $v = 101$       5 occurrences

Definition: A *finite word* is a finite sequence of letters belonging to a finite set called *alphabet*.

Example:  $101, 101001 \in \{0, 1\}^*$

## Binomial coefficient of words

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The *binomial coefficient*  $\binom{u}{v}$  of  $u$  and  $v$  is the number of times  $v$  occurs as a subsequence of  $u$  (meaning as a “scattered” subword).

Example:  $u = 101001$        $v = 101$       6 occurrences

Definition: A *finite word* is a finite sequence of letters belonging to a finite set called *alphabet*.

Example:  $101, 101001 \in \{0, 1\}^*$

## Binomial coefficient of words

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The *binomial coefficient*  $\binom{u}{v}$  of  $u$  and  $v$  is the number of times  $v$  occurs as a subsequence of  $u$  (meaning as a “scattered” subword).

Example:  $u = 101001$        $v = 101$

$$\Rightarrow \binom{101001}{101} = 6$$

Remark:

Natural generalization of binomial coefficients of integers

If  $a$  is a letter,

$$\binom{a^m}{a^k} = \binom{\overbrace{a \cdots a}^{m \text{ times}}}{\underbrace{a \cdots a}_{k \text{ times}}} = \binom{m}{k} \quad \forall m, k \in \mathbb{N}.$$



# Generalized Pascal triangles

Let  $(A, <)$  be a totally ordered alphabet.

Let  $L \subset A^*$  be an infinite language (set of words) over  $A$ .

The words in  $L$  are genealogically ordered

$$w_0 <_{\text{gen}} w_1 <_{\text{gen}} w_2 <_{\text{gen}} \dots$$

The *generalized Pascal triangle*  $P_L$  associated with  $L$  is defined by

$$P_L: (m, k) \in \mathbb{N} \times \mathbb{N} \mapsto \binom{w_m}{w_k} \in \mathbb{N}.$$

Questions:

- With a similar construction, can we expect the convergence to an analogue of the Sierpiński gasket?
- In particular, where should we cut to normalize a given generalized Pascal triangle?
- Could we describe this limit object?

Definitions:

- $\text{rep}_2(n)$  greedy base-2 representation of  $n \in \mathbb{N}_{>0}$  starting with 1
- $\text{rep}_2(0) = \varepsilon$  where  $\varepsilon$  is the empty word

$n$	$n = \sum_i c_i 2^i$ with $c_i \in \{0, 1\}$	$\text{rep}_2(n)$
0		$\varepsilon$
1	$1 \times 2^0$	1
2	$1 \times 2^1 + 0 \times 2^0$	10
3	$1 \times 2^1 + 1 \times 2^0$	11
4	$1 \times 2^2 + 0 \times 2^1 + 0 \times 2^0$	100
5	$1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0$	101
6	$1 \times 2^2 + 1 \times 2^1 + 0 \times 2^0$	110
$\vdots$	$\vdots$	$\vdots$
		$L_2 = 1\{0, 1\}^* \cup \{\varepsilon\}$

# Generalized Pascal triangle $P_2$ in base 2

$\binom{\text{rep}_2(m)}{\text{rep}_2(k)}$	$\text{rep}_2(k)$								
	$\varepsilon$	1	10	11	100	101	110	111	...
$\varepsilon$	1	0	0	0	0	0	0	0	
1	1	1	0	0	0	0	0	0	
10	1	1	1	0	0	0	0	0	
11	1	<b>2</b>	0	<b>1</b>	0	0	0	0	
$\text{rep}_2(m)$ 100	1	1	2	0	1	0	0	0	
101	1	2	1	1	0	1	0	0	
110	1	2	2	1	0	0	1	0	
111	1	3	0	<b>3</b>	0	0	0	1	
$\vdots$									$\ddots$

Rule (not local):

$$\binom{ua}{vb} = \binom{u}{vb} + \delta_{a,b} \binom{u}{v}$$

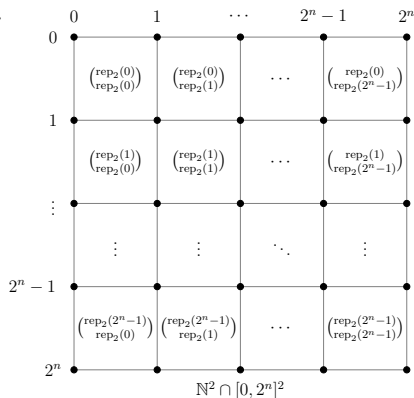
$(\text{rep}_2(m))$   
 $(\text{rep}_2(k))$  $\text{rep}_2(k)$ 

	$\epsilon$	<b>1</b>	10	<b>11</b>	100	101	110	<b>111</b>	$\dots$
$\epsilon$	<b>1</b>	0	0	0	0	0	0	0	
<b>1</b>	<b>1</b>	<b>1</b>	0	0	0	0	0	0	
10	1	1	1	0	0	0	0	0	
<b>11</b>	<b>1</b>	<b>2</b>	0	<b>1</b>	0	0	0	0	
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101	1	2	1	1	0	1	0	0	
110	1	2	2	1	0	0	1	0	
<b>111</b>	<b>1</b>	<b>3</b>	0	<b>3</b>	0	0	0	<b>1</b>	
$\vdots$									$\dots$

The classical Pascal triangle

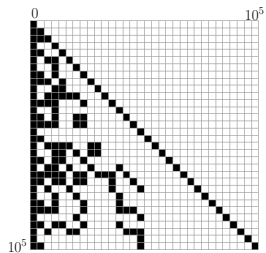
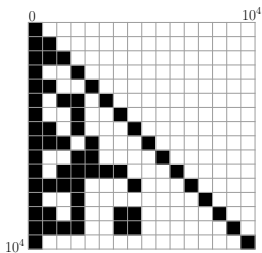
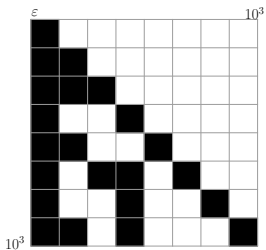
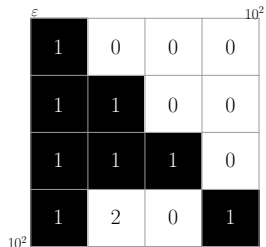
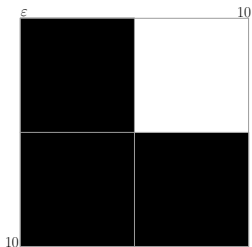
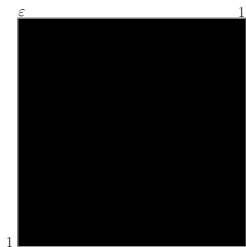
# Same construction

- Grid: first  $2^n$  rows and columns of  $\mathbb{P}_2$
- Color each square in
  - white if  $\binom{\text{rep}_2(m)}{\text{rep}_2(k)} \equiv 0 \pmod{2}$
  - black if  $\binom{\text{rep}_2(m)}{\text{rep}_2(k)} \equiv 1 \pmod{2}$
- Normalize by a homothety of ratio  $1/2^n$  (bring into  $[0, 1]^2$ )  
 $\rightsquigarrow$  sequence of compact sets belonging to  $[0, 1]^2$

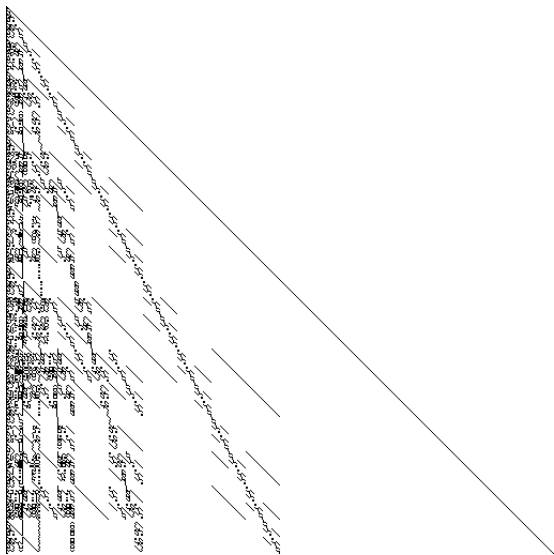


$$U_n = \frac{1}{2^n} \bigcup_{\substack{u, v \in L_2, |u|, |v| \leq n \\ \binom{u}{v} \equiv 1 \pmod{2}}} \text{val}_2(v, u) + [0, 1]^2$$

# The elements $U_0, \dots, U_5$



# The element $U_9$



Lines of different slopes: 1, 2, 4, 8, 16, ...

$(\star)$

$$(u, v) \text{ satisfies } (\star) \text{ iff } \begin{cases} u, v \neq \varepsilon \\ \binom{u}{v} \equiv 1 \pmod{2} \\ \binom{u}{v0} = 0 = \binom{u}{v1} \end{cases}$$

Example:  $(u, v) = (101, 11)$  satisfies  $(\star)$

$$\binom{101}{11} = 1$$

$$\binom{101}{110} = 0$$

$$\binom{101}{111} = 0$$



## Lemma: Completion

$(u, v)$  satisfies  $(\star) \Rightarrow (u_0, v_0), (u_1, v_1)$  satisfy  $(\star)$

Proof: Since  $(u, v)$  satisfies  $(\star)$

$$\begin{pmatrix} u \\ v \end{pmatrix} \equiv 1 \pmod{2}, \quad \begin{pmatrix} u \\ v_0 \end{pmatrix} = 0 = \begin{pmatrix} u \\ v_1 \end{pmatrix}$$

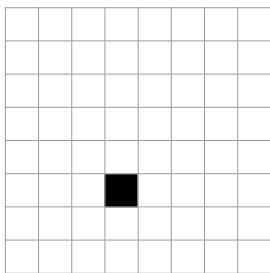
Proof for  $(u_0, v_0)$ :

$$\begin{pmatrix} u_0 \\ v_0 \end{pmatrix} = \underbrace{\begin{pmatrix} u \\ v_0 \end{pmatrix}}_{=0 \text{ since } (\star)} + \underbrace{\begin{pmatrix} u \\ v \end{pmatrix}}_{\equiv 1 \pmod{2}} \equiv 1 \pmod{2}$$

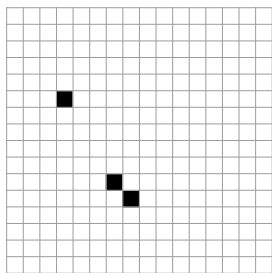
If  $\binom{u_0}{v_0} > 0$  or  $\binom{u_0}{v_0} > 0$ , then  $v_0$  is a subsequence of  $u$ .  
This contradicts  $(\star)$ .

Same proof for  $(u_1, v_1)$ . □

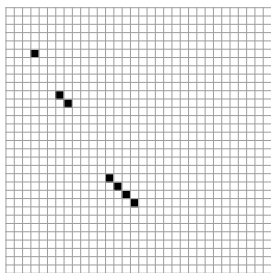
Example:  $(u, v) = (101, 11)$  satisfies  $(\star) \Rightarrow \binom{u}{v} \equiv 1 \pmod{2}$



$U_3$



$U_4$



$U_5$

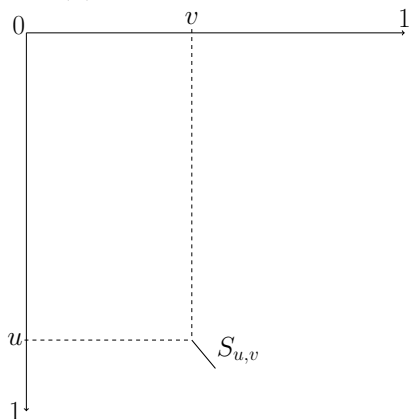
$\rightsquigarrow$  Creation of a segment of slope 1

Endpoint  $(3/8, 5/8) = (\text{val}_2(11)/2^3, \text{val}_2(101)/2^3)$

Length  $\sqrt{2} \cdot 2^{-3}$

# Segments of slope 1

The  $(\star)$  condition describes lines of slope 1 in  $[0, 1]^2$ .



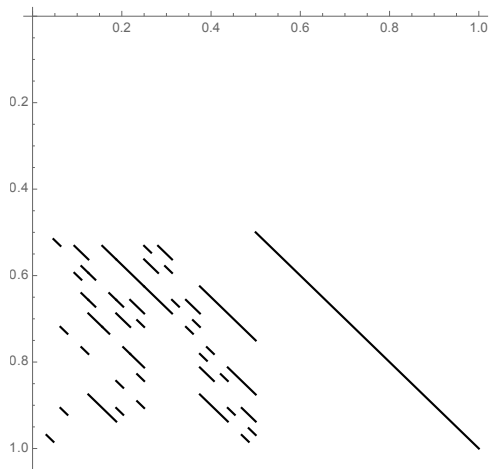
$(u, v) \in L_2 \times L_2$  satisfying  $(\star)$   
 $\rightsquigarrow$  closed segment  $S_{u,v}$

- slope 1
- length  $\sqrt{2} \cdot 2^{-|u|}$
- origin

$$\begin{aligned} A_{u,v} &= \text{val}_2(v, u) / 2^{|u|} \\ &= (0.0^{|u|-|v|}v, 0.u) \end{aligned}$$

Definition: New compact set containing those lines

$$\mathcal{A}_0 = \overline{\bigcup_{\substack{(u,v) \\ \text{satisfying } (\star)}} S_{u,v} \subset [0, 1]^2}$$

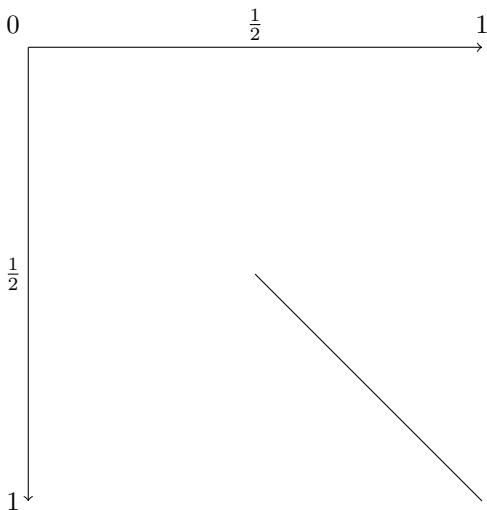


# Modifying the slope

Two maps  $c: (x, y) \mapsto (x/2, y/2)$  and  $h: (x, y) \mapsto (x, 2y)$

Example:  $(1, 1)$  satisfies  $(\star)$

Segment  $S_{1,1}$   
endpoint  $(1/2, 1/2)$   
length  $\sqrt{2} \cdot 2^{-1}$

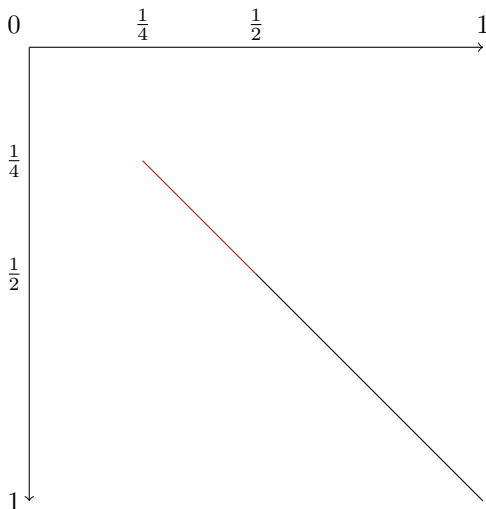


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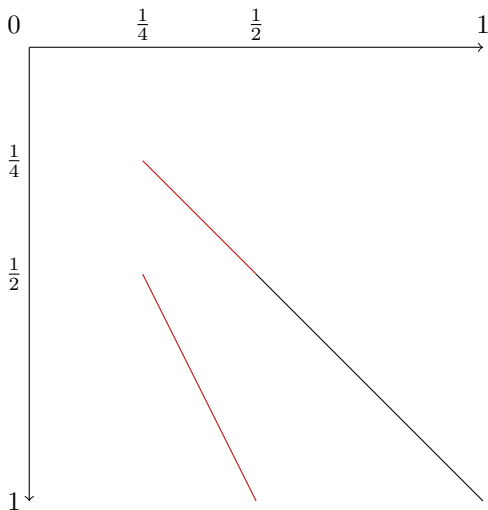


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endpoint  $(1/2, 1/2)$   
length  $\sqrt{2} \cdot 2^{-1}$

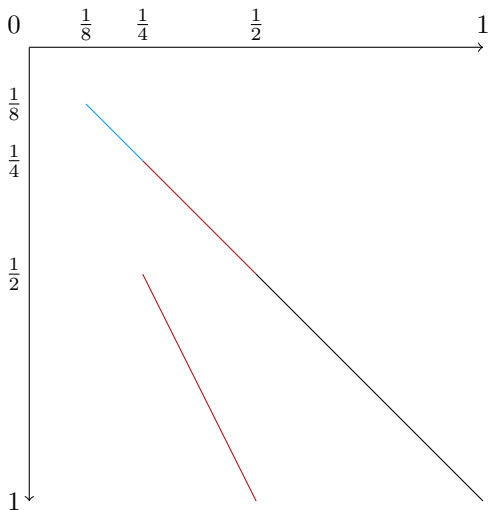


# Modifying the slope

Two maps  $c: (x, y) \mapsto (x/2, y/2)$  and  $h: (x, y) \mapsto (x, 2y)$

Example:  $(1, 1)$  satisfies  $(\star)$

Segment  $S_{1,1}$   
endpoint  $(1/2, 1/2)$   
length  $\sqrt{2} \cdot 2^{-1}$



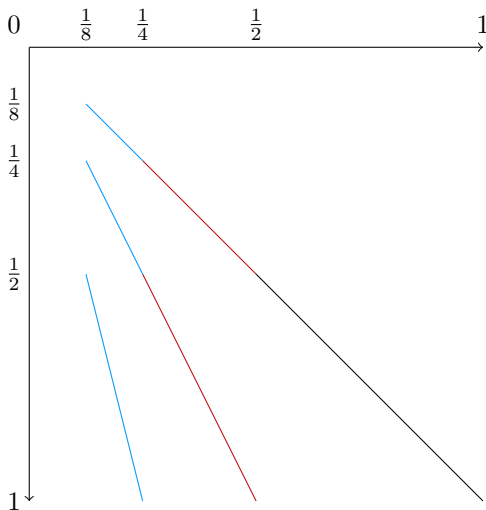


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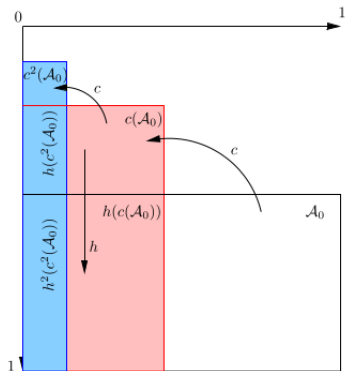


Definition: New compact set containing lines of slopes 1, 2,  $2^2, \dots, 2^n$

$$c: (x, y) \mapsto (x/2, y/2)$$

$$h: (x, y) \mapsto (x, 2y)$$

$$\mathcal{A}_n = \bigcup_{0 \leq j \leq i \leq n} h^j(c^i(\mathcal{A}_0))$$



The compact sets  $(\mathcal{A}_n)_{n \geq 0}$  are increasingly nested and their union is bounded.

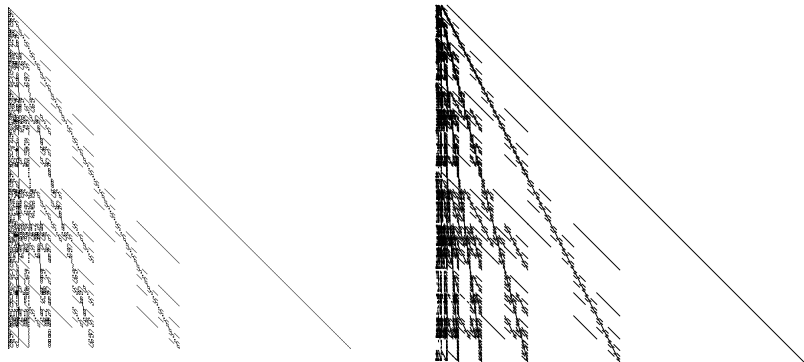
Thus  $(\mathcal{A}_n)_{n \geq 0}$  converges to

$$\mathcal{L} = \overline{\bigcup_{n \geq 0} \mathcal{A}_n}$$

(w.r.t. the Hausdorff distance).

## Theorem (Leroy, Rigo, S., 2016)

The sequence  $(U_n)_{n \geq 0}$  of compact sets converges to the compact set  $\mathcal{L}$  when  $n$  tends to infinity (w.r.t. the Hausdorff distance).



“Simple” characterization of  $\mathcal{L}$ :  $(\star)$  condition

Previous result: even and odd coefficients

## Theorem (Lucas, 1878)

Let  $p$  be a prime number.

If  $m = m_k p^k + \dots + m_1 p + m_0$  and  $n = n_k p^k + \dots + n_1 p + n_0$  then

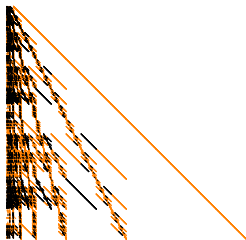
$$\binom{m}{n} \equiv \prod_{i=0}^k \binom{m_i}{n_i} \pmod{p}.$$

## Theorem (Leroy, Rigo, S., 2016)

Let  $p$  be a prime and  $0 < r < p$ .

When considering binomial coefficients congruent to  $r \pmod{p}$ , the sequence  $(U_{n,p,r})_{n \geq 0}$  converges to a well-defined compact set  $\mathcal{L}_{p,r}$  (w.r.t. the Hausdorff distance).

Example:  $\mathcal{L}_{3,1} \cup \mathcal{L}_{3,2}$

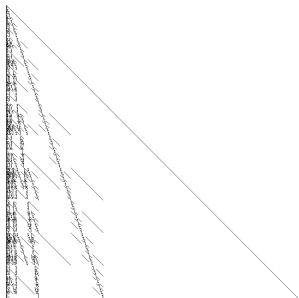


## Extension to any integer base

Everything still holds for binomial coefficients  $\equiv r \pmod{p}$  with

- integer base  $b \geq 2$
- language  $L_b$  of greedy base- $b$  representations of integers
- $p$  a prime
- $r \in \{1, \dots, p-1\}$

Example: base 3,  $\equiv 1 \pmod{2}$



# Fibonacci numeration system

## Definitions:

- Fibonacci numbers  $(F(n))_{n \geq 0}$   
 $F(0) = 1, F(1) = 2, F(n+2) = F(n+1) + F(n) \quad \forall n \geq 0$   
1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377, 610, 987, 1597, ...
- $\text{rep}_F(n)$  greedy Fibonacci representation of  $n \in \mathbb{N}_{>0}$  starting with 1
- $\text{rep}_F(0) = \varepsilon$  where  $\varepsilon$  is the empty word

$n$	$n = \sum_i c_i F(i)$ with $c_i \in \{0, 1\}$	$\text{rep}_F(n)$
0		$\varepsilon$
1	$1 \times F(0)$	1
2	$1 \times F(1) + 0 \times F(0)$	10
3	$1 \times F(2) + 0 \times F(1) + 0 \times F(0)$	100
4	$1 \times F(2) + 0 \times F(1) + 1 \times F(0)$	101
5	$1 \times F(3) + 0 \times F(2) + 0 \times F(1) + 0 \times F(0)$	1000
6	$1 \times F(3) + 0 \times F(2) + 0 \times F(1) + 1 \times F(0)$	1001
$\vdots$	$\vdots$	$\vdots$
		$L_F = 1\{0, 01\}^* \cup \{\varepsilon\}$

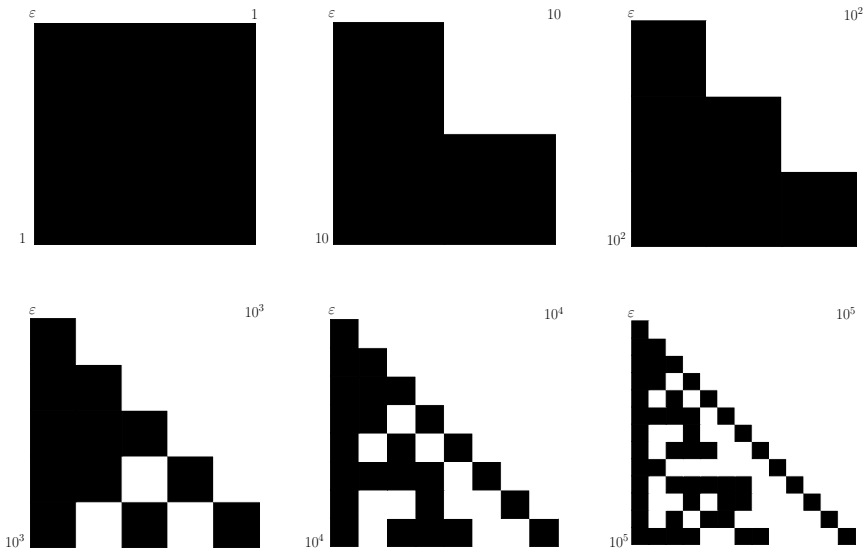
# Generalized Pascal triangle $P_F$ in Fibonacci base

$\begin{pmatrix} \text{rep}_F(m) \\ \text{rep}_F(k) \end{pmatrix}$	$\text{rep}_F(k)$								
	$\varepsilon$	1	10	100	101	1000	1001	1010	...
$\varepsilon$	1	0	0	0	0	0	0	0	
1	1	1	0	0	0	0	0	0	
10	1	1	1	0	0	0	0	0	
100	1	1	2	1	0	0	0	0	
$\text{rep}_F(m)$ 101	1	2	1	0	1	0	0	0	
1000	1	1	3	3	0	1	0	0	
1001	1	2	2	1	2	0	1	0	
1010	1	2	3	1	1	0	0	1	
$\vdots$									$\ddots$

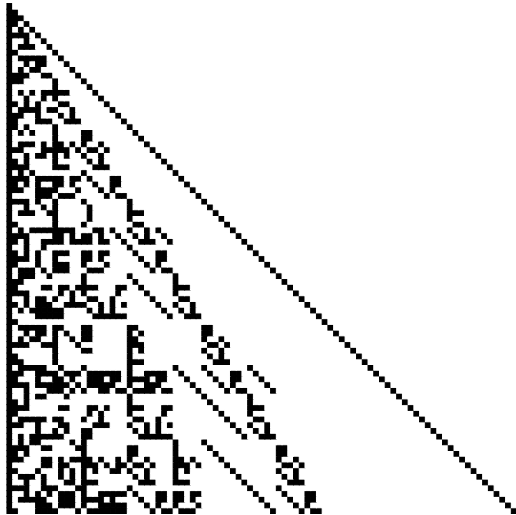
Rule (not local):

$$\begin{pmatrix} ua \\ vb \end{pmatrix} = \begin{pmatrix} u \\ vb \end{pmatrix} + \delta_{a,b} \begin{pmatrix} u \\ v \end{pmatrix}$$

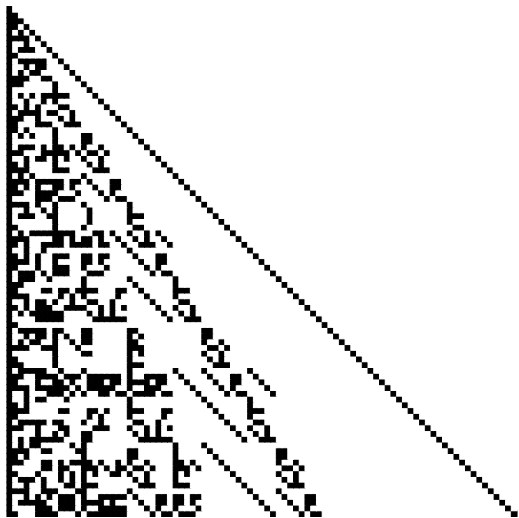
# The first six elements of the sequence $(U'_n)_{n \geq 0}$







Lines of different slopes:



Lines of different slopes:  $\varphi^n$ ,  $n \geq 0$ , with  $\varphi = \frac{1+\sqrt{5}}{2}$  Golden Ratio

Recall

- $(u, v)$  satisfies  $(\star)$  iff  $\binom{u}{v} \equiv 1 \pmod{2}$ , and  $\binom{u}{v0} = 0 = \binom{u}{v1}$ .
- $(u, v)$  satisfies  $(\star) \Rightarrow (u0, v0), (u1, v1)$  satisfy  $(\star)$

Problem: we cannot **always** add a letter 1 as a **suffix** in  $L_F$ .

Solution:  $p(u, v) \in \mathbb{N}$  s.t.  $u0^{p(u,v)}w, v0^{p(u,v)}w \in L_F$  for all  $w \in 0^*L_F$

$(\star')$

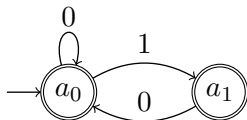
$$(u, v) \text{ satisfies } (\star') \text{ iff } u = v = \varepsilon \text{ or } \begin{cases} u, v \neq \varepsilon \\ \binom{u0^{p(u,v)}}{v0^{p(u,v)}} \equiv 1 \pmod{2} \\ \binom{u0^{p(u,v)}a}{v0^{p(u,v)}a} = 0 \quad \forall a \in \{0, 1\}. \end{cases}$$

# An automaton to find $p(u, v)$

Definition: A deterministic finite automaton (DFA) over an alphabet  $A$  is given by a 5-tuple  $\mathcal{A} = (Q, q_0, F, A, \delta)$  where

- $Q$  is a finite set of states
- $q_0 \in Q$  is the initial state (incoming arrow)
- $F \subset Q$  is the set of final states (concentric circles)
- $\delta: Q \times A \mapsto Q$  is the transition function

Example:



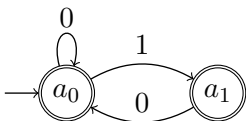
$$A = \{0, 1\}$$

$$Q = \{a_0, a_1\}$$

$$q_0 = a_0$$

$$F = \{a_0, a_1\}$$

$$\delta(a_0, 0) = a_0, \delta(a_0, 1) = a_1, \delta(a_1, 0) = a_0$$



Reading words in the automaton:

100101	$a_0 \xrightarrow{1} a_1 \xrightarrow{0} a_0 \xrightarrow{0} a_0 \xrightarrow{1} a_1 \xrightarrow{0} a_0 \xrightarrow{1} a_1$	✓	accepted
1011	$a_0 \xrightarrow{1} a_1 \xrightarrow{0} a_0 \xrightarrow{1} a_1 \xrightarrow{1} ???$	✗	not accepted
000101	$a_0 \xrightarrow{0} a_0 \xrightarrow{0} a_0 \xrightarrow{0} a_0 \xrightarrow{1} a_1 \xrightarrow{0} a_0 \xrightarrow{1} a_1$	✓	accepted

Accepted language (set of words):  $0^*L_F$  with  $L_F = 1\{0,01\}^* \cup \{\varepsilon\}$

Definition:  $p(u,v)$  is the smallest  $p$  s.t.  $\delta(a_0, u0^p) = a_0 = \delta(a_0, v0^p)$

Then  $u0^{p(u,v)}w, v0^{p(u,v)}w \in L_F$  for all  $w \in 0^*L_F$ .

- Completion lemma with the  $(\star')$  condition
- Creation of segments of slope 1
- New compact set  $\mathcal{A}'_0$  containing those lines

$$\mathcal{A}'_0 = \overline{\bigcup_{\substack{(u,v) \\ \text{satisfying } (\star')}} S_{u,v}} \subset [0, 1]^2$$

- Modification of the slopes with  $c: (x, y) \mapsto (x/\varphi, y/\varphi)$  and  $h: (x, y) \mapsto (x, \varphi y)$
- New compact set  $\mathcal{A}'_n$  containing lines of slopes  $1, \varphi, \varphi^2, \dots, \varphi^n$

$$\mathcal{A}'_n = \bigcup_{0 \leq j \leq i \leq n} h^j(c^i(\mathcal{A}'_0))$$

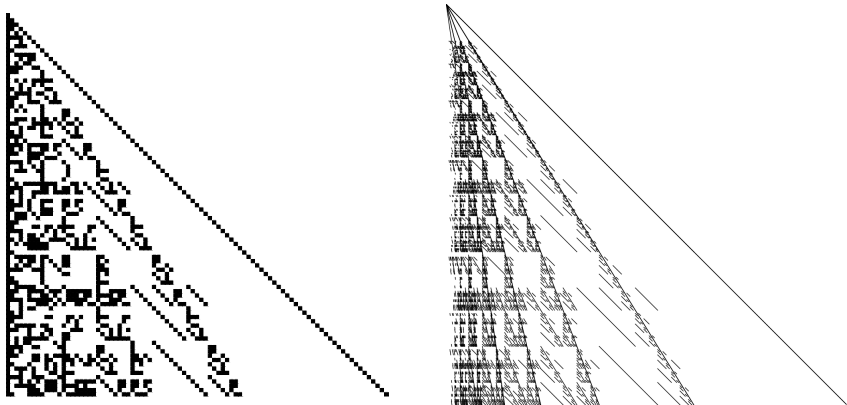
- $(\mathcal{A}'_n)_{n \geq 0}$  converges to

$$\mathcal{L}' = \overline{\bigcup_{n \geq 0} \mathcal{A}'_n}$$

(w.r.t. the Hausdorff distance).

## Theorem (S., 2018)

The sequence  $(U'_n)_{n \geq 0}$  of compact sets converges to the compact set  $\mathcal{L}'$  when  $n$  tends to infinity (w.r.t. the Hausdorff distance).



“Simple” characterization of  $\mathcal{L}'$ :  $(\star')$  condition

## Definition

A *numeration system* is a sequence  $U = (U(n))_{n \geq 0}$  of integers s.t.

- $U$  increasing
- $U(0) = 1$
- $\sup_{n \geq 0} \frac{U(n+1)}{U(n)}$  bounded by a constant  $\rightsquigarrow$  finite alphabet.

A numeration system  $U$  is *linear* if  $\exists k \geq 1, \exists a_0, \dots, a_{k-1} \in \mathbb{Z}$  s.t.

$$U(n+k) = a_{k-1}U(n+k-1) + \dots + a_0U(n) \quad \forall n \geq 0.$$

Greedy representation in  $(U(n))_{n \geq 0}$ :

$$n = \sum_{i=0}^{\ell} c_i U(i) \quad \text{with} \quad \sum_{i=0}^{j-1} c_i U(i) < U(j)$$

$$\text{rep}_U(n) = c_\ell \cdots c_0 \in \underbrace{L_U = \text{rep}_U(\mathbb{N})}_{\text{numeration language}}$$

Example: integer base  $(b^n)_{n \geq 0}$  with  $b \in \mathbb{N}_{>1}$

Fibonacci numeration system  $(F(n))_{n > 0}$



$$\beta \in \mathbb{R}_{>1} \quad A_\beta = \{0, 1, \dots, \lceil \beta \rceil - 1\}$$

$$x \in [0, 1] \rightsquigarrow x = \sum_{j=1}^{+\infty} c_j \beta^{-j}, \quad c_j \in A_\beta$$

Greedy way:  $c_j \beta^{-j} + c_{j+1} \beta^{-j-1} + \dots < \beta^{-(j-1)}$

$\beta$ -expansion of  $x$ :  $d_\beta(x) = c_1 c_2 c_3 \dots$

## Definition

$\beta \in \mathbb{R}_{>1}$  is a *Parry number* if  $d_\beta(1)$  is ultimately periodic.

Example:  $b \in \mathbb{N}_{>1}$ :  $d_b(1) = (b-1)^\omega$

Golden ratio  $\varphi$ :  $d_\varphi(1) = 110^\omega$

Parry number  $\beta \in \mathbb{R}_{>1} \rightsquigarrow$  linear numeration system  $(U_\beta(n))_{n \geq 0}$

- $d_\beta(1) = t_1 \cdots t_m 0^\omega$

$$\begin{aligned}U_\beta(0) &= 1 \\U_\beta(i) &= t_1 U_\beta(i-1) + \cdots + t_i U_\beta(0) + 1 & \forall 1 \leq i \leq m-1 \\U_\beta(n) &= t_1 U_\beta(n-1) + \cdots + t_m U_\beta(n-m) & \forall n \geq m\end{aligned}$$

- $d_\beta(1) = t_1 \cdots t_m (t_{m+1} \cdots t_{m+k})^\omega$

$$\begin{aligned}U_\beta(0) &= 1 \\U_\beta(i) &= t_1 U_\beta(i-1) + \cdots + t_i U_\beta(0) + 1 & \forall 1 \leq i \leq m+k-1 \\U_\beta(n) &= t_1 U_\beta(n-1) + \cdots + t_{m+k} U_\beta(n-m-k) & \forall n \geq m+k \\&+ U_\beta(n-k) \\&- t_1 U_\beta(n-k-1) - \cdots - t_m U_\beta(n-m-k)\end{aligned}$$

Examples:

$\bar{b} \in \mathbb{N}_{>1} \rightsquigarrow (b^n)_{n \geq 0}$  base  $b$

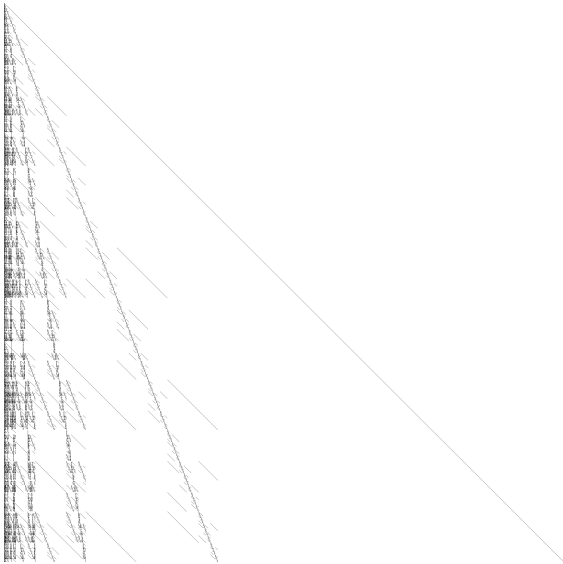
Golden ratio  $\varphi \rightsquigarrow (F(n))_{n \geq 0}$  Fibonacci numeration system

- Parry number  $\beta \in \mathbb{R}_{>1}$
- Parry numeration system  $(U_\beta(n))_{n \geq 0}$
- Numeration language  $L_{U_\beta}$
- Generalized Pascal triangle  $P_\beta$  in  $(U_\beta(n))_{n \geq 0}$  indexed by words of  $L_{U_\beta}$
- Sequence of compact sets extracted from  $P_\beta$  (first  $U_\beta(n)$  rows and columns of  $P_\beta$ )
- Convergence to a limit object (same technique)
  - Lines of different slopes:  $\beta^n, n \geq 0$
  - $(\star')$  condition and description of segments of slope 1
  - Two maps  $c: (x, y) \mapsto (x/\beta, y/\beta)$  and  $h: (x, y) \mapsto (x, \beta y)$
  - Sequence of sets  $\mathcal{A}_n^\beta$  containing lines of slopes  $1, \beta, \beta^2, \dots, \beta^n$
  - $\mathcal{A}_n^\beta$  converges to

$$\mathcal{L}^\beta = \overline{\bigcup_{n \geq 0} \mathcal{A}_n^\beta}$$

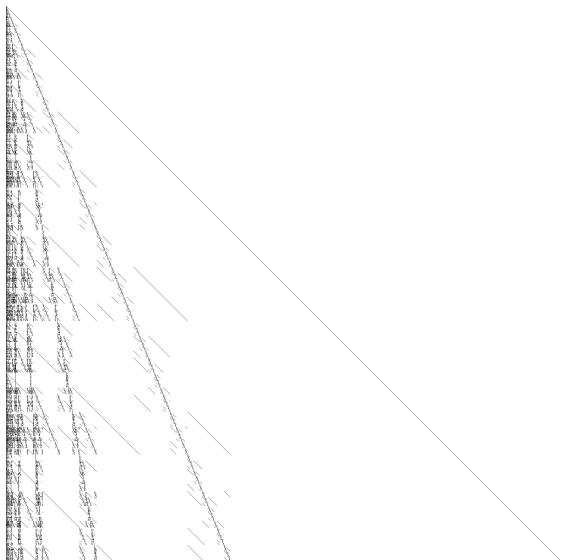
- Works modulo any prime number

# Example 1

 $\varphi^2$ 

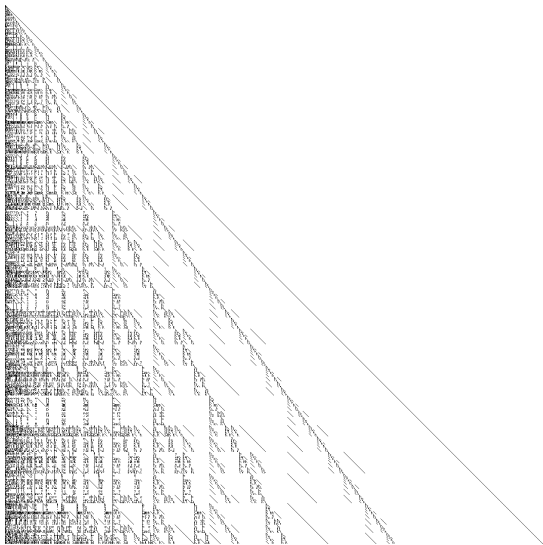
## Example 2

$\beta_1 \approx 2.47098$  dominant root of  $P(X) = X^4 - 2X^3 - X^2 - 1$



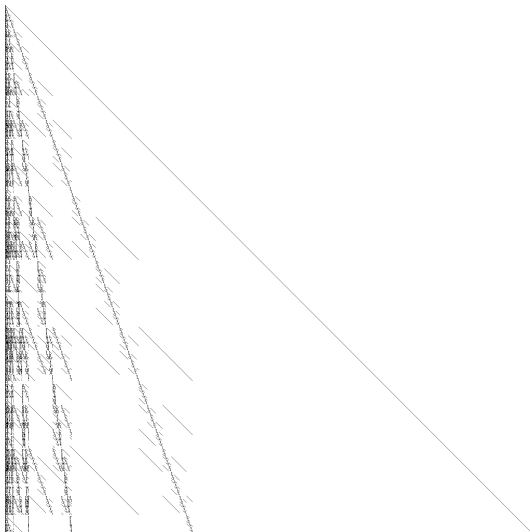
## Example 3

$\beta_2 \approx 1.38028$  dominant root of  $P(X) = X^4 - X^3 - 1$



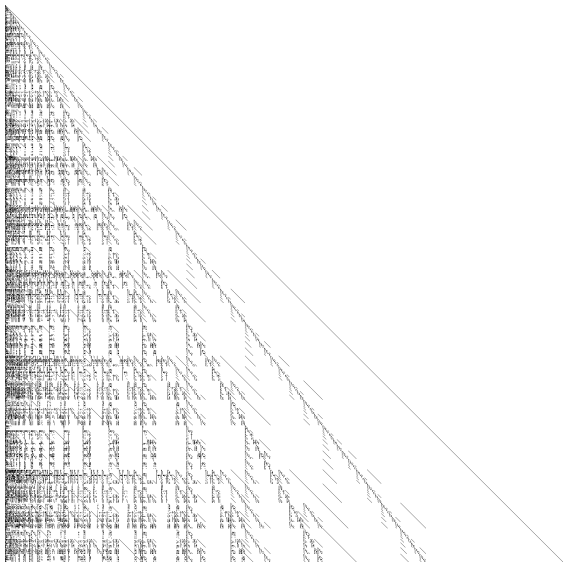
## Example 4

$\beta_3 \approx 2.80399$  dominant root of  $P(X) = X^4 - 2X^3 - 2X^2 - 2$



## Example 5

$\beta_4 \approx 1.32472$  dominant root of polynomial  $P(X) = X^5 - X^4 - 1$





In this talk:

Numeration system	Generalized Pascal triangle	Convergence mod $p$
Base 2	✓	✓
Integer base	✓	✓
Fibonacci	✓	✓
Parry	✓	✓

- Regularity of the sequence counting subword occurrences: result for any integer base  $b$  and the Fibonacci numeration system
- Behavior of the summatory function: result for any integer base  $b$  (exact behavior) and the Fibonacci numeration system (asymptotics)

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