Nyldon Words

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The lexicographic order

Let A be a finite alphabet with a total order <.

Then the order < on the letters induces a total order $<_{lex}$ on A^* called the lexicographic order:

$$x <_{\text{lex}} y \iff x \in \operatorname{Pref}(y) \setminus y$$

or $\exists p \in A^* \& a, b \in A$:
$$\begin{cases} a < b \\ pa \in \operatorname{Pref}(x) \\ pb \in \operatorname{Pref}(y) \end{cases}$$

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Lyndon words

A finite word w over A is Lyndon if for all x, y such that w = xy, we have $w <_{lex} yx$.

Thus, w is Lyndon if and only if it is primitive and minimal among its conjugates.

- ▶ 0,1
- ▶ 01

▶ ...

- ▶ 001,011
- ▶ 0001,0011,0111
- 00001,00011,00101,00111,01011,01111

The Chen-Fox-Lyndon theorem

Every finite word w over A can be uniquely factorized as

$$w = \ell_1 \ell_2 \cdots \ell_k$$

where $k \in \mathbb{N}$ and $(\ell_i)_{1 \le i \le k}$ is a nonincreasing sequence of Lyndon words.

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Some Lyndon factorizations:

- $\bullet 0100010110 = (01)(0001011)(0)$
- > 000100111001 = (000100111001)
- 0110101 = (011)(01)(01)

Recursive definition of Lyndon words

- The letters are Lyndon.
- A finite word is Lyndon if and only if it cannot be factorized as a nonincreasing sequence of shorter Lyndon words.

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Nyldon words

- The letters are Nyldon.
- A finite word is Nyldon if and only if it cannot be factorized as a nondecreasing sequence of shorter Nyldon words.

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Nyldon words were defined in a post on Mathoverflow by Grinberg in November 2014, together with 2 conjectures:

- Is it true that any finite word can be uniquely factorized as a nondecreasing sequence of shorter Nyldon words?
- Is it true that Nyldon words form a set of representatives of the primitive conjugacy classes?

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▶ 00,01,10,11



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- ▶ 00, 01, 10, 11
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 100,101,110,111

0,110

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 100,101,110,111

0,1
10
100,101

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 100,101,110,111
- 0000,0001,0010,0011,
 0100,0101,0110,0111,
 1000,1001,1010,1011,
 1100,1101,1110,1111

0,1
10
100,101

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- 000,001,010,011,
 100,101,110,111
- 0000,0001,0010,0011,
 0100,0101,0110,0111,
 1000,1001,1010,1011,
 1100,1101,1110,1111

- ▶ 0,1
- ► 10
- ▶ 100,101
- ▶ 1000, 1001, 1011

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- ▶ 0,1
- ► 10

▶ ...

- ▶ 100, 101
- ▶ 1000, 1001, 1011
- 10000, 10001, 10010, 10011, 10110, 10111
- 100000, 100001, 100010, 100011, 100110, 100111, 101100, 101110, 101111
- 1000000, 1000001, 1000010, 1000011, 1000100, 1000110, 1000111, 1001100, 1001110, 1001111, 1011000, 1011001, 1011100, 1011101, 1011110, 1011111

If we count them, we find: $2, 1, 2, 3, 6, 9, 18, 30, 56, 99, 186, \ldots$

First observations and properties

- Nyldon words are not maximal among their conjugates: for example, 101 is Nyldon.
- Nyldon words are not extremal among their conjugates, for any order we have thought of.
- Apart from 0 and 1, all binary Nyldon words start with the prefix 10.
- If one can prove the unicity of the Nyldon factorization, then we know that they are as many Nyldon words as Lyndon words of each length.

Nyldon words start with the prefix 10: why is that?

Nyldon words can't start with the letter 0 (except for 0 itself).

• Let
$$w = 0u$$
, with $u \neq \varepsilon$.

Write

$$u = n_1 \cdots n_k, \qquad k \ge 1, \quad n_1 \le_{\text{lex}} \cdots \le_{\text{lex}} n_k$$

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- But $0 \leq_{\text{lex}} n_1$ and 0 is Nyldon.
- ► Thus w = (0)(n₁) ··· (n_k) is a factorization of w into nondecreasing shorter Nyldon words.
- By definition, *w* is not Nyldon.

Nyldon words start with the prefix 10: why is that?

Nyldon words can't start with 11.

• Let
$$w = 11u$$
.

Write

$$1u = n_1 \cdots n_k, \qquad k \ge 1, \quad n_1 \le_{\text{lex}} \cdots \le_{\text{lex}} n_k$$

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- Now observe that n_1 begins in 1.
- Then $1 \leq_{\text{lex}} n_1$ and 1 is Nyldon.
- ► This shows that w = 11u = (1)(n₁) · · · (n_k) factors into nondecreasing shorter Nyldon words.
- By definition, *w* is not Nyldon.

Forbidden prefixes

We can extend this argument in order to show that many other prefixes are not allowed in the family of Nyldon words:

▶ 00,01

▶ ...

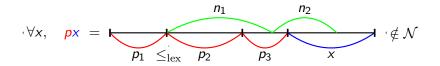
- ▶ 11, 1010, 100100, ... (in general $10^k 10^k$)
- ▶ 1001011, 10001011, . . . (in general 10^k1011)
- ▶ 1011011, 101110111,... (in general 101^k01^k)
- ▶ 10011011, 1000110011, ... (in general $10^{k+1}110^{k}11$)

Forbidden prefixes

A forbidden prefix p is a finite word such that no Nyldon word starts with p.

All the examples we have are from the following family:

$$F = \{ p \in A^* \colon p = p_1 p_2 p_3, \ p_1 \text{ Nyldon}, \ p_1 \leq_{\text{lex}} p_2, \\ \forall x \in A^*, \ p_2 p_3 x = n_1 \cdots n_k \implies |n_1| \ge |p_2| \}$$



- ► All elements of *F* are forbidden prefixes.
- ► Question: Are there other forbidden prefixes?

Suffixes, rather than prefixes

- Prefixes of Lyndon and Nyldon words behave in a very different manner.
- On the opposite, we can show that suffixes of both families share some properties.

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Theorem (Lyndon)

Let $w \in A^*$. Then the following assertions are equivalent:

- w is Lyndon
- w is strictly smaller than all its proper suffixes
- w is strictly smaller than all its proper Lyndon suffixes.

Theorem (Nyldon)

A Nyldon word is strictly greater than all its proper Nyldon suffixes.

But the condition is not sufficient in the Nyldon case, even if we ask it to be true for all proper suffixes: 110 is not Nyldon.

Nyldon suffixes of Nyldon words

Nyldon

Lyndon

1000000, 1000001, 1000010, 1000011, 1000100, 1000110, 1000111, 1001100, 1001110, 1001111, 1011000, 1011001, 1011100, 1011101, 1011110, 1011111

0000001,0000011,0000101, 0000111,0001001,0001101, 0001111,0010011,0011101, 0011111,0001011,0011011, 0010111,0110111,0101111, 0111111

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The longest Nyldon suffix of a word

Theorem (Lyndon)

- The longest Lyndon suffix of a finite word is the right-most factor of its Lyndon factorization.
- The longest Lyndon prefix of a finite word is the left-most factor of its Lyndon factorization.

Theorem (Nyldon)

The longest Nyldon suffix of a finite word is the right-most factor of any of its Nyldon factorizations.

NB: There is no similar condition on prefixes for Nyldon words.

Unicity of the Nyldon factorization

Every finite word w over A can be uniquely factorized as

 $w = n_1 n_2 \cdots n_k$

where $k \in \mathbb{N}$ and $(n_i)_{1 \le i \le k}$ is a nondecreasing sequence of Nyldon words.

Lyndon vs Nyldon factorizations:

- 0100010110 = (01)(0001011)(0) = (0)(1000)(10110)
- $\bullet 000100111001 = (000100111001) = (0)(0)(0)(100111001)$
- 0110101 = (011)(01)(01) = (0)(1)(10)(101)

Standard factorization

Theorem (Lyndon)

- Let S be the longest Lyndon proper suffix of a word w and let w = PS. Then w is Lyndon iff P is Lyndon and $P <_{lex} S$.
- Let P be the longest Lyndon proper prefix of a word w and let w = PS. Then w is Lyndon iff S is Lyndon and $P <_{\text{lex}} S$.

Theorem (Nyldon)

Let S be the longest Nyldon proper suffix of a word w and let w = PS. Then w is Nyldon iff P is Nyldon and $P >_{lex} S$.

NB: There is no similar condition on prefixes for Nyldon words.

Faster algorithm for computing Nyldon words

Compute the Nyldon factorization of w from right to left.

```
n \leftarrow |w|: NyIF \leftarrow (w[n])
for i = 1 to n - 1 do
   if w[n-i] \leq_{\text{lex}} N_{Y}/F(1) then
      NyIF \leftarrow (w[n-i], NyIF)
   else
      while \#Ny|F \ge 2 and Ny|F(1) >_{\text{lex}} Ny|F(2) do
         NyIF \leftarrow (NyIF(1) \cdot NyIF(2), NyIF(3), \ldots)
      end while
   end if
end for
return NylF
```

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• w is Nyldon if and only if \#NyIF = 1.
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From Lyndon to Nyldon: lost properties

Many results for Lyndon words don't have analogues in the case of Nyldon words.

- Results concerning prefixes.
- ▶ Let u, v be Lyndon words such that u <_{lex} v. Then uv is Lyndon.
- ▶ For example, 10010 >_{lex} 1 but 100101 = (100)(101) is not Nyldon.
- The longest Lyndon suffix of a word is also the suffix that is minimal w.r.t. the lexicographic order.
- For example, the longest Nyldon suffix of 110 is 10 although it is not lex-maximal (nor lex-minimal).

Complete factorizations of the free monoid

A totally ordered family $F \subseteq A^*$ is called a complete factorization of A^* if each $w \in A^*$ can be uniquely factorized as

$$w = x_1 x_2 \cdots x_k \tag{1}$$

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where $k \in \mathbb{N}$ and $(x_i)_{1 \le i \le k}$ is a nonincreasing sequence of F.

- ► Lyndon words are a complete factorization for the order <_{lex}.
- ▶ Nyldon words are a complete factorization for the order >_{lex}.

Schützenberger's theorem (1965)

Quote from the webpage of Dominique Perrin: "The following result has no known elementary proof."

Theorem (Schützenberger 1965)

Let A be an alphabet, $F \subseteq A^*$ and \prec be a total order on F. Then any two of the following three conditions imply the third:

- Each $w \in A^*$ admits at least one factorization (1).
- Each $w \in A^*$ admits at most one factorization (1).
- All elements of F are primitive and each primitive conjugacy class of A⁺ contains exactly one element of F.

As a consequence of the unicity of the Nyldon factorization and Schützenberger's theorem, we obtain:

Corollary

The Nyldon words form a set of representatives of the primitive conjugacy classes.

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A simpler proof of the primitivity of Nyldon words?

Lyndon words are primitive by definition.

But suppose for a minute that you only know Lyndon words from their recursive definition and that all Lyndon proper suffixes of a Lyndon word is greater than the word itself.

(Of course, also suppose that you don't know Schützenberger's theorem.)

Then we easily deduce that Lyndon word are primitive and that each primitive conjugacy class contains exactly one Lyndon word. Lyndon word are primitive and each primitive conjugacy class contains exactly one Lyndon word

- By induction on the length *n* of the words.
- Base case: all letters are Lyndon.
- Let w be a word of length n ≥ 2 which is a power: w = x^m with x primitive and m ≥ 2.
- By induction hypothesis, we know that x has a Lyndon conjugate: y = vu is Lyndon and x = uv.

• Then
$$w = x^m = (uv)^m = u(vu)^{m-1}v = uy^m v$$
.

• Let $u = u_1 \cdots u_k$ and $v = v_1 \cdots v_\ell$ be the Lyndon factorizations of u and v.

• Then
$$u_k \geq_{\text{lex}} y \geq_{\text{lex}} v_1$$
.

• Therefore $w = u_1 \cdots u_k \cdot \underbrace{y \cdots y}_{l} \cdot v_1 \cdots v_\ell$ is not Lyndon.

Lyndon word are primitive and each primitive conjugacy class contains exactly one Lyndon word.

- We have shown that Lyndon words of length *n* are primitive.
- Now suppose that there exist distinct Lyndon words x, y of length n in the same conjugacy class: x = uv and y = vu.
- Then x^2 has two Lyndon factorizations:

$$x^{2} = x \cdot x$$
$$= u_{1} \cdots u_{k} \cdot y \cdot v_{1} \cdots v_{\ell}.$$

 But our hypotheses imply the unicity of the Lyndon factorization, hence we have reached a contradiction. The same reasoning doesn't work for Nyldon words.

If x = uv and y = vu is Nyldon, then

$$x^m = u_1 \cdots u_k \cdot \underbrace{y \cdots y}_{m-1} \cdot v_1 \cdots v_\ell$$

is not the Nyldon factorization of x^m .

• For example, if x = 01 then y = 10 and $x^2 = (0)(101)$.

If yv is Nyldon, then the Nyldon factorization of x^m is

$$x^m = u_1 \cdots u_k \cdot \underbrace{y \cdots y}_{m-2} \cdot (yv).$$

But this is not always the case.

► For example, if x = 001100101 then y = 100101001 and x² = (0)(0)(1)(10010)(1001100)(101).

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Lazard factorizations of the free monoid

A Lazard factorization of A^* is a set $F \subseteq A^*$ satisfying the following conditions:

- 1. F is totally ordered by some order \prec
- 2. For every $n \ge 1$, if

$$F \cap A^{\leq n} = \{u_1, \ldots, u_k\}$$
 with $u_1 \prec \cdots \prec u_k$

and if

$$Y_1 = A$$
 and $Y_{i+1} = u_i^*(Y_i \setminus u_i)$ for all i,

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then we have

(i) for every
$$i \in \{1, \dots, k\}$$
, $u_i \in Y_i$
(ii) $Y_k \cap A^{\leq n} = \{u_k\}$.

Lyndon is Lazard

The set \mathcal{L} is a Lazard factorization for the order $<_{lex}$. For $A = \{0, 1\}$, the words of length ≤ 4 in the lexicographic order are 0,0001,001,0011,011,0111,1.

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We observe that $u_i = \min Y_i$.

Nyldon is not Lazard

Theorem (Viennot 1978)

A complete factorization (F, \prec) of A^{*} is a Lazard factorization if and only if

$$\forall x, y \in F, \, xy \in F \implies x \prec xy.$$

- Since the Lyndon words are a complete factorization w.r.t. <_{lex}, this result confirms that they are indeed a Lazard factorization.
- Since the Nyldon words are a complete factorization w.r.t.
 >lex, this result tells us that they are *not* a Lazard factorization.

But Nyldon is right-Lazard

For $A = \{0, 1\}$, the Nyldon words of length ≤ 4 in the (increasing) lexicographic order are 0, 1, 10, 100, 1000, 1001, 101, 1011.

We observe that $u_i = \min Y_i$.

Right Lazard factorizations of the free monoid

Lazard factorization \rightarrow left Lazard factorization

A right Lazard factorization of A^* is a set $F \subseteq A^*$ satisfying the following conditions:

- 1. F is totally ordered by some order \prec
- 2. For every $n \ge 1$, if

$$F \cap A^{\leq n} = \{u_1, \ldots, u_k\}$$
 with $u_1 \succ \cdots \succ u_k$

and if

$$Y_1 = A$$
 and $Y_{i+1} = (Y_i \setminus u_i) u_i^*$ for all i,

then we have

(i) for every
$$i \in \{1, \ldots, k\}$$
, $u_i \in Y_i$
(ii) $Y_k \cap A^{\leq n} = \{u_k\}$.

Regular factorization = left and right Lazard factorization. =

Right factorizations of the free monoid

Theorem (Viennot 1978)

A complete factorization (F, \prec) of A^* is a right Lazard factorization if and only if

$$\forall x, y \in F, \, xy \in F \implies xy \prec y.$$

- \blacktriangleright Lyndon words are a right Lazard factorization for $<_{\rm lex}.$
- ▶ Nyldon words are a right Lazard factorization for >_{lex}.

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Four types of Lazard factorizations

The min and max choices show an asymetric situation for left and right Lazard factorizations:

Left Lazard		Right Lazard	
lexmin	L	lexmin	\mathcal{N}
lexmax	$\bar{\mathcal{L}}$	lexmax	${\cal L}$

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One can now easily check that $\overline{\mathcal{L}}$ is a complete factorization of A^* for the order \propto defined by $u \propto v$ if and only if $\overline{u} <_{lex} \overline{v}$.

This emphasizes that a complete factorization of the free monoid is given by the choice of the order on the words, more than by the choice of the allowed factors.

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