# Permutations and shifts: a survey 

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## Permutations

We write

$$
312=\left(\begin{array}{lll}
1 & 2 & 3 \\
3 & 1 & 2
\end{array}\right)
$$

and

$$
(312)=\left(\begin{array}{lll}
1 & 2 & 3 \\
2 & 3 & 1
\end{array}\right)=231
$$

## Let's start with the tent map

$$
T:[0,1] \rightarrow[0,1], x \mapsto\left\{\begin{array}{cl}
2 x & \text { if } x \in\left[0, \frac{1}{2}\right] \\
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## Let's start with the tent map



A case study shows


So the permutation $\pi=321$ is not realizable!

## Dynamical systems

- $(X, T)$ where $X$ is a set and $T$ is a map from $X$ to $X$.
- The objects of study are the trajectories of points.
- The orbit of $x \in X$ is the subset $\left\{T^{n}(x): n \in \mathbb{N}\right\}$.
- Typically, the set $X$ is endowed with a specific structure and the map $T$ preserves this structure.
- If $X$ is a topological space and $T$ is continuous, then $(X, T)$ is a topological dynamical system.
- If $X$ is a measurable space and $T$ is measure preserving, then $(X, T)$ is a measure-preserving dynamical system.


## Conjugacy (in the topological case)

- $\left(X_{1}, T_{1}\right)$ and $\left(X_{2}, T_{2}\right)$ are conjugate if there exists a homeomorphism $\phi: X_{1} \rightarrow X_{2}$ such that $\phi \circ T_{1}=T_{2} \circ \phi$

$$
\begin{array}{cl}
X_{1} \xrightarrow{T_{1}} X_{1} \\
\phi \downarrow \\
\downarrow \\
X_{2} \xrightarrow{T_{2}}{ }^{\downarrow}{ }^{\downarrow}{ }_{2}
\end{array}
$$

- One of the goals in the theory: classify dynamical systems up to conjugacy.


## Invariants

- The idea: if two conjugate systems necessarily share some property, which is called an invariant, then this property can be used to distinguish non-conjugate systems.
- The useful invariants must be computable for a large class of dynamical systems.
- Example of invariant: the number of periodic points.


## Entropy

- Permits us to measure the complexity of a dynamical system.
- Invariant under conjugacy.
- Computable for a large variety of dynamical systems.
- So, it is a powerful tool in order to classify dynamical systems.


## The starting point

Let $I$ be an interval of $\mathbb{R}$ and consider the dynamical systems $(I, T)$ where $T: I \rightarrow I$.

Theorem (Bandt-Keller-Pompe 2002)

- The concepts of permutation entropy and of topological entropy coincide for piecewise monotone interval maps.
- Similar result for the Kolmogorov-Sinai entropy w.r.t. an invariant measure.

Entropy of interval maps via permutations [Bandt-Keller-Pompe 2002]

## Permutation entropy

- Let $(X, T)$ be a dynamical system where $X$ is a totally ordered set.
- For an integer $n \geq 1$ and a point $x \in X$ such that

$$
x, T(x), \ldots, T^{n-1}(x)
$$

are pairwise distinct, $\operatorname{Pat}(T, n, x)$ denotes the permutation $\pi \in \mathcal{S}_{n}$ defined by

$$
T^{\pi^{-1}(1)-1}(x)<T^{\pi^{-1}(2)-1}(x)<\cdots<T^{\pi^{-1}(n)-1}(x) .
$$

- Otherwise stated, $\pi(i)<\pi(j)$ for all $i, j \in \llbracket 1, n \rrbracket$ such that $T^{i-1}(x)<T^{j-1}(x)$.

Example
If $T^{3}(x)<T(x)<x<T^{2}(x)$ then $\operatorname{Pat}(T, 4, x)=3241$.

## Permutation entropy

- $\operatorname{Allow}(T, n)=\{\operatorname{Pat}(T, n, x): x \in X\}$ is the set of permutations of length $n$ realized by some $x \in X$.
- $\operatorname{Allow}(T)=\bigcup_{n \geq 1} \operatorname{Allow}(T, n)$.
- The permutation entropy of $T$ is defined as

$$
\lim _{n \rightarrow \infty} \frac{1}{n} \log |\operatorname{Allow}(T, n)|
$$

provided that the limit exists.

## Symbolic dynamical systems

- The idea is to discretize dynamical systems.
- The set $X$ is partitioned into subsets $P_{1}, \ldots, P_{k}$.
- A point $x \in X$ is coded by a right-infinite word $\left(a_{n}\right)_{n \in \mathbb{N}}$ :

$$
\forall n \in \mathbb{N}, a_{n}=i \text { whenever } T^{n}(x) \in P_{i}
$$

- If a point $x \in X$ is coded by $\left(a_{n}\right)_{n \in \mathbb{N}}$, then its image $T(x)$ is coded by $\left(a_{n+1}\right)_{n \in \mathbb{N}}$.

- We are interested in determining which sequences can arise in this way.


## Binary representation of numbers

$$
\text { Let } T:[0,1) \rightarrow[0,1), x \mapsto\{2 x\}
$$



We partition $[0,1)$ into the 2 subintervals $\left[0, \frac{1}{2}\right)$ and $\left[\frac{1}{2}, 1\right)$, which are coded by 0 and 1 respectively.
Then the coding of a real number $x$ just corresponds to its binary expansion.

## Representation of numbers in a real base $\beta$

Let $\beta>1$ a real number and $T_{\beta}:[0,1) \rightarrow[0,1), x \mapsto\{\beta x\}$.


$$
\beta=\sqrt{5}
$$

We partition $[0,1)$ into the $\lceil\beta\rceil$ subintervals

$$
\left[0, \frac{1}{\beta}\right),\left[\frac{1}{\beta}, \frac{2}{\beta}\right), \ldots,\left[\frac{\lceil\beta\rceil-1}{\beta}, 1\right),
$$

which are coded by $0,1, \ldots,\lceil\beta\rceil-1$ respectively.
In this case the coding of a real number $x$ corresponds to its $\beta$-expansion.

## Symbolic dynamical systems

- $X$ is a subset of $A^{\mathbb{N}}$ stable under the shift operator $\sigma$ :

$$
\sigma\left(\left(a_{n}\right)_{n \in \mathbb{N}}\right)=\left(a_{n+1}\right)_{n \in \mathbb{N}}
$$

- Denote by $\sigma_{X}$ the restriction of the operator $\sigma$ to $X$.
- If $X$ is also compact then $\left(X, \sigma_{X}\right)$ is called a symbolic dynamical system, a shift space, or simply a shift.
- A shift can also be described as a set $X_{\mathcal{F}}$ of all sequences avoiding the finite blocks in $\mathcal{F}$.
- $\left(A^{\mathbb{N}}, \sigma\right)$ is called the full shift.


## Entropy in shifts $\left(X, \sigma_{X}\right)$

- $\operatorname{Fact}_{n}(X)$ is the number of factors of length $n$ that appear in some $x \in X$.
- Typically, $\left|\operatorname{Fact}_{n}(X)\right|$ grows like $2^{c n}$ for some constant $c$.
- The entropy of $\sigma_{X}$ is given by

$$
\lim _{n \rightarrow \infty} \frac{1}{n} \log \left|\operatorname{Fact}_{n}(X)\right|
$$

- We can equip $A^{\mathbb{N}}$, and hence any shift space, with a total order, as the lexicographic order for example.
- Using Bandt-Keller-Pompe's result, an alternative way to compute the entropy is given by

$$
\lim _{n \rightarrow \infty} \frac{1}{n} \log \left|\operatorname{Allow}\left(\sigma_{X}, n\right)\right|
$$

## Uncountably many forbidden permutations

- The same result implies that not all permutations are realizable in such a dynamical system.
- In fact, in general, there are much more forbidden permutations than realizable permutations.
- Quote from Elizalde: "Understanding the forbidden patterns of chaotic maps is important because the absence of these patterns is what distinguishes sequences generated by chaotic maps from random sequences."


## Part I: Permutations in full shifts

Forbidden patterns and shift systems [Amigó-Elizalde-Kennel 2008] The number of permutations realized by a shift [Elizalde 2009]

## The full shift over $k$ symbols

- Let $A_{k}=\{0,1, \ldots, k-1\}$ and $\sigma_{k}: A_{k}^{\mathbb{N}} \rightarrow A_{k}^{\mathbb{N}}$ denote the shift operator.
- Elements in $A_{k}^{\mathbb{N}}$ are ordered by the lexicographic order:

$$
\begin{array}{r}
a_{1} a_{2} \cdots<_{\text {lex }} b_{1} b_{2} \cdots \Longleftrightarrow \exists i \geq 1, a_{1} \cdots a_{i-1}=b_{1} \cdots b_{i-1} \\
\text { and } a_{i}<b_{i}
\end{array}
$$

- Study the permutations realizable in full shifts $\left(A_{k}^{\mathbb{N}}, \sigma_{k}\right)$, that is, the sets $\operatorname{Allow}\left(\sigma_{k}\right)$.
- In particular, for a given permutation $\pi$, compute the quantity

$$
N_{+}(\pi)=\min \left\{k \geq 1: \pi \in \operatorname{Allow}\left(\sigma_{k}\right)\right\}
$$

which is the number of symbols needed in order to realize $\pi$.

Example ( $\pi=4217536 \in \mathcal{S}_{7}$ )
Then $\operatorname{Pat}\left(\sigma_{3}, 7,210221220 \cdots\right)=\pi$ since

| $210221220 \cdots$ | 4 |
| :--- | :--- |
| $10221220 \cdots$ | 2 |
| $0221220 \cdots$ | 1 |
| $221220 \cdots$ | 7 |
| $21220 \cdots$ | 5 |
| $1220 \cdots$ | 3 |
| $220 \cdots$ | 6 |

Another way to see it is:

$$
\begin{aligned}
& 210221220 \ldots \\
& 4217536
\end{aligned}
$$

In fact, to realize the permutation $\pi$, one needs 3 symbols, so that $N_{+}(\pi)=3$.

## Computing $N_{+}(\pi)$

- Associated with $\pi \in \mathcal{S}_{n}$, we consider the circular permutation (or n-cycle)

$$
\hat{\pi}=(\pi(1) \pi(2) \cdots \pi(n)),
$$

that is, $\hat{\pi}(\pi(i))=\pi(i+1)$ for $1 \leq i<n$, and $\hat{\pi}(\pi(n))=\pi(1)$.

- Idea: Count the number of descents in $\hat{\pi}$.
- A descent in a permutation $\pi \in \mathcal{S}_{n}$ is an index $1 \leq i<n$ such that $\pi(i)>\pi(i+1)$.


## Example

If we represent the permutation $\pi=2413$, we see that it has one descent:


## Computing $N_{+}(\pi)$

Theorem (Elizalde 2009)
For any $\pi \in \mathcal{S}_{n}$, the minimal number $k$ of distinct symbols of a sequence $w$ satisfying $\operatorname{Pat}\left(w, \sigma_{k}, n\right)=\pi$ is

$$
N_{+}(\pi)=1+\operatorname{des}(\hat{\pi})+\epsilon_{+}(\pi)
$$

where $\operatorname{des}(\hat{\pi})$ is the number of descents in $\hat{\pi}$ with $\pi(1)$ removed, and

$$
\epsilon_{+}(\pi)= \begin{cases}1 & \text { if } \pi \text { ends with } 21 \text { or with }(n-1) n \\ 0 & \text { otherwise }\end{cases}
$$

## Computing $N_{+}(\pi)$ : Sketch of the proof

- If $\left\{\begin{array}{l}a_{i} a_{i+1} \cdots<_{\text {lex }} a_{j} a_{j+1} \cdots \\ a_{i}=a_{j}\end{array}\right.$ then $a_{i+1} a_{i+2} \cdots<_{\text {lex }} a_{j+1} a_{j+2} \cdots$
- Now suppose that $a_{1} a_{2} \cdots \in A_{k}^{\mathbb{N}}$ realizes the permutation $\pi \in \mathcal{S}_{n}$, that is, $\operatorname{Pat}\left(\sigma_{k}, n, a_{1} a_{2} \cdots\right)=\pi$.
- If $\left\{\begin{array}{l}\pi(i)<\pi(j) \\ a_{i}=a_{j} \\ i, j<n\end{array} \quad\right.$ then $\pi(i+1)<\pi(j+1)$.
- We make use of $\hat{\pi}$ with the contrapositive statement.
- If $\left\{\begin{array}{l}\pi(i)+1=\pi(j) \\ \pi(i+1)=\hat{\pi}(\pi(i))>\hat{\pi}(\pi(j))=\pi(j+1) \quad \text { then } a_{i}<a_{j} \text {. } \\ i, j<n\end{array}\right.$
- So, for each descent in $\hat{\pi}$ with $\pi(1)$ removed, we need one more symbol.

Example $\left(\pi=4217536 \in \mathcal{S}_{7}\right)$
One has $\hat{\pi}=(4217536)=7162345$ and $\operatorname{des}(\hat{\pi})=2$.

$$
\hat{\pi}=7 \begin{array}{llllll}
7 & 1 & 6 & 2 & 3 & \underline{4}
\end{array}
$$

Finding digits

$$
\pi=\begin{array}{lllllll}
4 & 2 & 1 & 7 & 5 & 3 & 6
\end{array}
$$

Placing digits

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|  | $\hat{\pi}=$ | 7 | 1 | 6 | 2 | 3 | $\underline{4}$ | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 1 |  |  |  |  |  |  |
| Finding digits |  |  |  |  |  |  | 1 | 7 |
|  | 5 | 3 | 6 |  |  |  |  |  |

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| Finding digits |  | 1 | 1 | 2 | 2 |  | 2 |  |
|  | $\pi=$ | 2 | 1 | 7 | 5 | 3 | 6 |  |
|  |  |  |  |  |  |  |  |  |
| Placing digits |  | 1 | 0 | 2 | 2 | 1 |  |  |

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One has $\hat{\pi}=(4217536)=7162345$ and $\operatorname{des}(\hat{\pi})=2$.

|  | $\hat{\pi}=$ | 7 | 1 | 6 | 2 | 3 | 4 | 5 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Finding digits |  | 1 | 1 | 2 | 2 |  | 2 |  |
|  | $\pi=$ | 2 | 1 | 7 | 5 | 3 | 6 |  |
| Placing digits |  | 1 | 0 | 2 | 2 | 1 |  |  |

If you ask for at most 3 symbols, then the prefix of any sequence realizing $\pi$ starts with the prefix $z_{1} \cdots z_{n-1}=210221$.

Example $\left(\pi=4217536 \in \mathcal{S}_{7}\right)$
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We can continue this prefix (using only 3 symbols) to obtain a sequence that realizes $\pi$.

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| Placing digits |  | 1 | 0 | 2 | 2 | 1 | 2 | 2 | 0 | $\ldots$ |  |

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| Placing digits |  | 1 | 0 | 2 | 2 | 1 | 2 | 2 | 0 | $\ldots$ |  |

If you ask for at most 3 symbols, then the prefix of any sequence realizing $\pi$ starts with the prefix $z_{1} \cdots z_{n-1}=210221$.

We can continue this prefix (using only 3 symbols) to obtain a sequence that realizes $\pi$.

In fact any sequence starting with the prefix 210221220 works.

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| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
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We can continue this prefix (using only 3 symbols) to obtain a sequence that realizes $\pi$.

In fact any sequence starting with the prefix 210221220 works.

## Computing $N_{+}(\pi)$

Theorem (Elizalde 2009)
For any $\pi \in \mathcal{S}_{n}$, the minimal number $k$ of distinct symbols of a sequence $w$ satisfying $\operatorname{Pat}\left(w, \sigma_{k}, n\right)=\pi$ is

$$
N_{+}(\pi)=1+\operatorname{des}(\hat{\pi})+\epsilon_{+}(\pi)
$$

where $\operatorname{des}(\hat{\pi})$ is the number of descents in $\hat{\pi}$ with $\pi(1)$ removed and

$$
\epsilon_{+}(\pi)= \begin{cases}1 & \text { if } \pi \text { ends with } 21 \text { or with }(n-1) n \\ 0 & \text { otherwise } .\end{cases}
$$

## Permutations in full shifts

- The shortest forbidden permutations of $A_{k}^{\mathbb{N}}$, have length $k+2$.
- For every $\pi \in \mathcal{S}_{n}$ we have $N_{+}(\pi) \leq n-1$.
- There are exactly 6 permutations $\pi$ in $\mathcal{S}_{n}$ such that $N_{+}(\pi)=n-1:$

$$
\begin{array}{ll}
1 n 2(n-1) 3(n-2) \ldots, & \ldots(n-2) 3(n-1) 2 n 1, \\
n 1(n-1) 2(n-2) 3 \ldots, & \ldots 3(n-2) 2(n-1) 1 n, \\
\ldots 4(n-1) 3 n 21, & \ldots(n-3) 2(n-2) 1(n-1) n .
\end{array}
$$

## Permutations in full shifts

- In fact, Elizalde shows much more by proving a closed formula for the number $a_{n, \ell}$ of permutations $\pi$ of length $n$ for which $N_{+}(\pi)=\ell$, for any $n$ and $\ell$.
- In particular, for each fixed $\ell, a_{n, \ell} \sim n \ell^{n-1}$ as $n \rightarrow \infty$.
- Then, for each $k$,

$$
\lim _{n \rightarrow \infty} \frac{1}{n} \log \left|\operatorname{Allow}\left(\sigma_{k}, n\right)\right|=\lim _{n \rightarrow \infty} \frac{1}{n} \log \left(\sum_{\ell=1}^{k} a_{n, \ell}\right)=\log k,
$$

in accordance with Bandt-Keller-Pompe's theorem.

## Part II: Permutations in $\beta$-shifts

Permutations and $\beta$-shifts [Elizalde 2011]

## Permutations in $\beta$-shifts

- For $\beta>1$ we study the dynamical systems $\left([0,1), T_{\beta}\right)$ where $T_{\beta}:[0,1) \rightarrow[0,1), x \mapsto\{\beta x\}$.

$\beta=2$

$\beta=\sqrt{5}$
- Study the realizable/forbidden permutations.


## The $\beta$-shift

- Instead of numbers $x \in[0,1)$, we will rather consider their $\beta$-expansions, denoted by $d_{\beta}(x)$.
- We let $\Omega_{\beta}$ denote the topological closure of the set $\left\{d_{\beta}(x): x \in[0,1)\right\}$ and $\sigma_{\beta}: \Omega_{\beta} \rightarrow \Omega_{\beta},\left(a_{m}\right) \mapsto\left(a_{m+1}\right)$.
- The map $\sigma_{\beta}$ is continuous and $\Omega_{\beta}$ is a compact metric space, hence the $\beta$-shift $\left(\Omega_{\beta}, \sigma_{\beta}\right)$ is a topological dynamical system.
- The case $\beta \in \mathbb{N}$ corresponds to full shifts.


## $\operatorname{Allow}\left(T_{\beta}\right)=\operatorname{Allow}\left(\sigma_{\beta}\right)$

- Key observation: $x<y \Longleftrightarrow d_{\beta}(x)<_{\text {lex }} d_{\beta}(y)$.
- The following diagram commutes

$$
\begin{gathered}
{[0,1) \xrightarrow{T_{\beta}}[0,1)} \\
d_{\beta} \left\lvert\, \begin{array}{l}
\| \\
\\
\Omega_{\beta} \xrightarrow{\sigma_{\beta}} \xrightarrow{\mid d_{\beta}} \\
\Omega_{\beta}
\end{array}\right.
\end{gathered}
$$

- Thus, $\boldsymbol{T}_{\beta}$ and $\sigma_{\beta}$ are order-isomorphic, and, for all $x \in[0,1)$ and all $n \geq 1$, we have

$$
\operatorname{Pat}\left(T_{\beta}, n, x\right)=\operatorname{Pat}\left(\sigma_{\beta}, n, d_{\beta}(x)\right)
$$

with the lexicographic order on $\Omega_{\beta}$.

## The shift complexity

- If $1<\beta \leq \beta^{\prime}$ then $\Omega_{\beta} \subseteq \Omega_{\beta^{\prime}}$ and $\operatorname{Allow}\left(T_{\beta}\right) \subseteq \operatorname{Allow}\left(T_{\beta^{\prime}}\right)$.
- Compute $B_{+}(\pi)=\inf \left\{\beta>1: \pi \in \operatorname{Allow}\left(T_{\beta}\right)\right\}$.
- This quantity is called the (positive) shift complexity of $\pi$.


## Example

For $n=2$, one has $B_{+}(12)=B_{+}(21)=1$.
For $n=3$, one has

$$
B_{+}(132)=B_{+}(213)=B_{+}(321)=\frac{1+\sqrt{5}}{2}
$$

and

$$
B_{+}(123)=B_{+}(231)=B_{+}(312)=1 .
$$

## Computing the shift complexity

- For $a=a_{1} a_{2} \cdots$ such that $a=\sup a_{j} a_{j+1} \cdots \neq \overline{0}$, let $b_{+}(a)$ be the unique solution $\beta \geq 1$ of

$$
\sum_{j=1}^{\infty} \frac{a_{j}}{\beta^{j}}=1
$$

By convention, $b_{+}(\overline{0})=1$.

- If $a$ is eventually periodic then $b_{+}(a)$ is the unique real root greater than or equal to 1 of a polynomial.


## Computing the shift complexity

- For $\pi \in \mathcal{S}_{n}$, define $z_{1} z_{2} \cdots z_{n-1}$ as in the case of full shifts.
- Let $m=\pi^{-1}(n)$ and $\ell=\pi^{-1}(\pi(n)-1)$ if $\pi(n) \neq 1$.

Theorem (Elizalde 2011)
Let $\pi \in \mathcal{S}_{n}$ and $\beta>1$. Then $\pi \in \operatorname{Allow}\left(T_{\beta}\right) \Longleftrightarrow \beta>b_{+}(a)$ where

$$
a= \begin{cases}z_{[m, n)} \overline{z_{[\ell, n)}} & \text { if } \pi(n) \neq 1, \\ z_{[m, n)} \overline{0} & \text { if } \pi(n)=1 \text { and } \pi(n-1) \neq 2 \\ z_{[m, n)}^{\prime} \overline{0} & \text { if } \pi(n)=1 \text { and } \pi(n-1)=2\end{cases}
$$

where for $1 \leq j<n, z_{j}^{\prime}=z_{j}+1$. In particular, $B_{+}(\pi)=b_{+}(a)$.
Theorem (Elizalde 2011)
We always have $\pi \notin \operatorname{Allow}\left(T_{B_{+}(\pi)}\right)$. So $N_{+}(\pi)=1+\left\lfloor B_{+}(\pi)\right\rfloor$.

## Minimal shift complexity

The only permutations $\pi \in \mathcal{S}_{n}$ satisfying $B_{+}(\pi)=1$ are

$$
(c+1)(c+2) \ldots n 12 \ldots c
$$

for any fixed $1 \leq c \leq n$.

Example
We already saw that

$$
B_{+}(12)=B_{+}(21)=1
$$

and

$$
B_{+}(123)=B_{+}(231)=B_{+}(312)=1 .
$$

## Maximal shift complexity

For $n=3$, there is 3 permutations of maximal complexity.

Theorem (Elizalde 2011)
For $\pi \in \mathcal{S}_{n} \backslash\left\{\rho_{n}\right\}$ with $n \geq 4$, we have $B_{+}(\pi)<B_{+}\left(\rho_{n}\right)$ where

$$
\rho_{n}= \begin{cases}1 n 2(n-1) \ldots \frac{n}{2} \frac{n+2}{2} & \text { if } n \text { is even } \\ 1 n 2(n-1) \ldots \frac{n-1}{2} \frac{n+3}{2} \frac{n+1}{2} & \text { if } n \text { is odd. }\end{cases}
$$

Moreover, $B_{+}\left(\rho_{n}\right) \in[n-2, n-1)$.

Example
We have $\rho_{4}=1423$ and $B_{+}\left(\rho_{4}\right)=\frac{3+\sqrt{5}}{2}=2.61 \ldots$

## $B_{+}\left(\rho_{n}\right)$ is the threshold

Recall that $\pi \notin$ Allow $\left(T_{B_{+}(\pi)}\right)$. Therefore we get
Corollary
For $n \geq 4$, we have $\mathcal{S}_{n} \subseteq \operatorname{Allow}\left(T_{\beta}\right) \Longleftrightarrow \beta>B_{+}\left(\rho_{n}\right)$.

Example (continued)
For $\beta>2.61 \ldots$, the $\beta$-shift allows all permutations of length $\leq 4$.

Corollary
For a fixed $\beta>1$, the length of the shortest forbidden permutation of $T_{\beta}$ is the integer $n \geq 2$ defined by $B_{+}\left(\rho_{n-1}\right)<\beta \leq B_{+}\left(\rho_{n}\right)$.

## Part III: Permutations and negative $\beta$-shifts

Patterns of negative shifts and beta-shifts [Elizalde-Moore] Permutations and negative beta-shifts [Charlier-Steiner]

## Negative $\beta$-shifts

- Let $\beta>1$. We study the map

$$
T_{-\beta}:(0,1] \rightarrow(0,1], x \mapsto\lfloor\beta x\rfloor+1-\beta x .
$$

- Generalization of $T_{\beta}$ as $T_{-\beta}(x)=\{-\beta x\}$ except for finitely many points.


$$
\beta=2
$$


$\beta=\frac{3+\sqrt{5}}{2}$

## Negative $\beta$-shifts

- Again, instead of numbers $x \in(0,1]$, we consider their $(-\beta)$-expansions, denoted by $d_{-\beta}(x)$.
- $\Omega_{-\beta}$ is the closure of $\left\{d_{-\beta}(x): x \in(0,1]\right\}$.
- The shift map is $\sigma_{-\beta}: \Omega_{-\beta} \rightarrow \Omega_{-\beta},\left(a_{m}\right) \mapsto\left(a_{m+1}\right)$.


## Permutations in negative $\beta$-shifts

- Key observation: $x<y \Longleftrightarrow d_{-\beta}(x)<_{\text {alt }} d_{-\beta}(y)$.
- Here we use the alternating lexicographic order for sequences:

$$
\begin{array}{r}
a_{1} a_{2} \cdots<_{\text {alt }} b_{1} b_{2} \cdots \Longleftrightarrow \exists i \geq 1, a_{1} \cdots a_{i-1}=b_{1} \cdots b_{i-1} \\
\\
\text { and } \begin{cases}a_{i}<b_{i} & \text { if } i \text { is odd }, \\
a_{i}>b_{i} & \text { if } i \text { is even. }\end{cases}
\end{array}
$$

For example, $1320 \cdots<_{\text {alt }} 1210 \cdots<_{\text {alt }} 1220 \cdots$

- We have $\operatorname{Allow}\left(T_{-\beta}\right)=\operatorname{Allow}\left(\sigma_{-\beta}\right)$ with the alternating lexicographic order on the $(-\beta)$-shift.


## Count the number of ascents

- We have to adapt the arguments from the full shift case.
- We consider again $\hat{\pi}=(\pi(1) \pi(2) \cdots \pi(n))$.
- Idea: For each ascent in $\hat{\pi}$ with $\pi(1)$ removed, we need one more symbol.



## Theorem (Charlier-Steiner, Elizalde-Moore)

Let $\pi \in \mathcal{S}_{n}$. Then the minimal number of symbols of a sequence $w$ satisfying $\operatorname{Pat}\left(\sigma_{k}, n, w\right)=\pi$ w.r.t. the alternating lexicographic order is

$$
N_{-}(\pi)=1+\operatorname{asc}(\hat{\pi})+\epsilon_{-}(\pi),
$$

where $\operatorname{asc}(\hat{\pi})$ is the number of ascents in $\hat{\pi}$ with $\pi(1)$ removed and

$$
\epsilon_{-}(\pi)= \begin{cases}1 & \text { if some condition on } \pi \text { holds } \\ 0 & \text { otherwise }\end{cases}
$$

In particular $N_{-}(\pi) \leq n-1$ for all $\pi \in \mathcal{S}_{n}, n \geq 3$.
For $n \geq 4$, there are exactly 4 permutations $\pi \in \mathcal{S}_{n}$ with $N_{-}(\pi)=n-1$ :
$12 \ldots n, \quad 12 \ldots(n-2) n(n-1), \quad n(n-1) \ldots 1, \quad n(n-1) \ldots 312$.

## Permutations in negative $\beta$-shifts

- Study the permutations realizable/forbidden in negative $\beta$-shifts.
- Compute the negative shift complexity

$$
B_{-}(\pi)=\inf \left\{\beta>1: \pi \in \operatorname{Allow}\left(T_{-\beta}\right)\right\}
$$

## Computing the negative shift complexity

Theorem (Charlier-Steiner, Elizalde-Moore)
Let $\pi \in \mathcal{S}_{n}$ and $\beta>1$. Then $\pi \in \operatorname{Allow}\left(T_{-\beta}\right) \Longleftrightarrow \beta>b_{-}(a)$ where
$a= \begin{cases}z_{[m, n)} \overline{z_{l \ell, n)}} & \text { if } n-m \text { is even, } \pi(n) \neq 1, \text { and }(\star), \\ \min _{0 \leq i<|r-\ell|} z_{[m, n)}^{(i)} \overline{z_{[\ell, n)}^{(i)}} & \text { if } n-m \text { is even, } \pi(n) \neq 1, \text { and } \neg(\star), \\ \overline{z_{[m, n)}^{0}} & \text { if } n-m \text { is even and } \pi(n)=1, \\ z_{[m, n)} \overline{z_{[r, n)}} & \text { if } n-m \text { is odd and } \neg(\star), \\ \min _{0 \leq i<|r-\ell|} z_{[m, n)}^{(i)} \overline{z_{[r, n)}^{(i)}} & \text { if } n-m \text { is odd and }(\star) .\end{cases}$
In particular $B_{-}(\pi)=b_{-}(a)$.

Theorem (Charlier-Steiner, Elizalde-Moore)
We have $N_{-}(\pi)=1+\left\lfloor B_{-}(\pi)\right\rfloor$.

## Minimal negative shift complexity

Theorem (Charlier-Steiner)
If $a \gg_{\text {alt }} \varphi^{\omega}(0)$ where $\varphi: 0 \mapsto 1,1 \mapsto 100$, then $B_{-}(\pi)$ is a Perron number, i.e., an algebraic integer $\beta>1$ all of whose Galois conjugates $\alpha$ satisfy $|\alpha|<\beta$.

Moreover, $B_{-}(\pi)=1 \Longleftrightarrow a=\overline{\varphi^{k}(0)}$ for some $k \geq 0$.

## Comparing the positive and negative $\beta$-shifts

| $B_{ \pm}(\pi)$ | root of | $\pi$, negative beta-shift | $\pi$, positive beta-shift |
| :---: | :---: | :---: | :---: |
| 1 | $\beta-1$ | 12,21 | 12,21 |
|  |  | $123,132,213,231,321$ | $123,231,312$ |
|  |  | $1324,1342,1432,2134$ | $1234,2341,3412,4123$ |
|  |  | $2143,2314,2431,3142$ |  |
| 1.465 | $\beta^{3}-\beta^{2}-1$ |  | $1342,2413,3124,4231,3421,4213$ |
| 1.618 | $\beta^{2}-\beta-1$ | $1423,3412,4231$ | $1243,1324,2431,3142,4312$ |
| 1.755 | $\beta^{3}-2 \beta^{2}+\beta-1$ | $2341,2413,3124,4123$ |  |
| 1.802 | $\beta^{3}-2 \beta^{2}-2 \beta+1$ |  | 4213 |
| 1.839 | $\beta^{3}-\beta^{2}-\beta-1$ | 4132 | $1432,2143,3214,4321$ |
| 2 | $\beta-2$ | 1234,1243 | 2134,3241 |
| 2.247 | $\beta^{3}-2 \beta^{2}-\beta+1$ | 4321 | 4132 |
| 2.414 | $\beta^{2}-2 \beta-1$ |  | 2314,3421 |
| 2.618 | $\beta^{2}-3 \beta+1$ |  | 1423 |
| 2.732 | $\beta^{2}-2 \beta-2$ | 4312 |  |

## Some open problems

- Count all permutations with $B_{-}(\pi) \leq N$ or $B_{-}(\pi)<N$, in particular with $B_{-}(\pi)=1$. From Bandt-Keller-Pompe's theorem we know that

$$
\lim _{n \rightarrow \infty} \frac{1}{n} \log \#\left\{\pi \in \mathcal{S}_{n}: B_{-}(\pi) \leq \beta\right\}=\log \beta
$$

What are the precise asymptotics of

$$
c_{n}=\#\left\{\pi \in \mathcal{S}_{n}: B_{-}(\pi)=1\right\} ?
$$

We have $\left(c_{n}\right)_{n \geq 2}=2,5,12,19,34,57,82,115, \ldots$

- Describe the permutations given by the transformations

$$
T_{\beta, \alpha}:[0,1) \rightarrow[0,1), x \mapsto \beta x+\alpha-\lfloor\beta x+\alpha\rfloor .
$$

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